



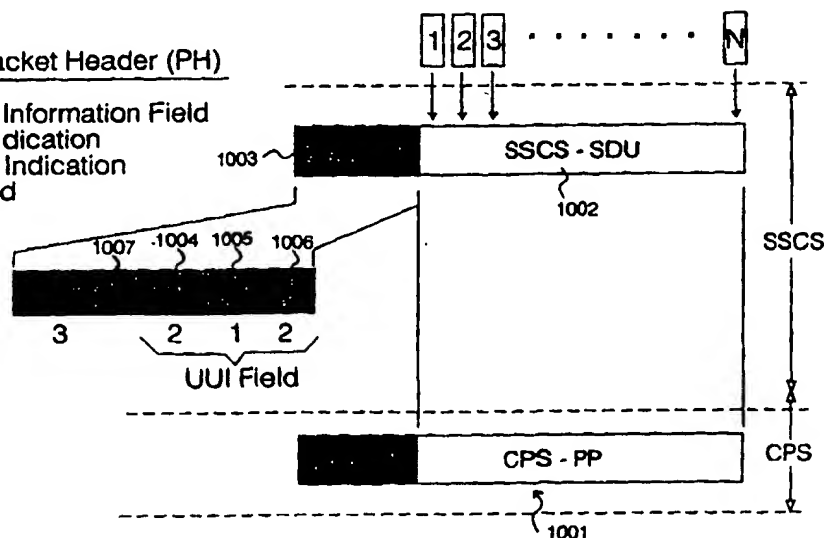
INTERNATIONAL APPLICATION PUBLISHED UNDER THE PATENT COOPERATION TREATY (PCT)

(51) International Patent Classification ⁶ : H04Q 11/04	A1	(11) International Publication Number: WO 98/33355 (43) International Publication Date: 30 July 1998 (30.07.98)
(21) International Application Number: PCT/GB98/00179 (22) International Filing Date: 21 January 1998 (21.01.98) (30) Priority Data: 9701270.2 22 January 1997 (22.01.97) GB (71) Applicant (for all designated States except US): NORTHERN TELECOM LIMITED [CA/CA]; World Trade Center of Montreal, 8th floor, 380 St. Antoine Street West, Montreal, Quebec H2Y 3Y4 (CA). (72) Inventors; and (75) Inventors/Applicants (for US only): STACEY, David, John [GB/GB]; 28 Orchard Close, Stanstead Abbots, Herts SG12 8AH (GB). BRUECKHEIMER, Simon, Daniel [GB/GB]; 74 Church Crescent, London N10 3NE (GB). CAVES, Keith [GB/GB]; 28 Sheering Lower Road, Sawbridgeworth, Herts CM21 9LF (GB). (74) Agent: RYAN, John, Peter, William; Nortel Patents, London Road, Harlow, Essex CM17 9NA (GB).		(81) Designated States: CA, JP, US, European patent (AT, BE, CH, DE, DK, ES, FI, FR, GB, GR, IE, IT, LU, MC, NL, PT, SE). Published <i>With international search report.</i> <i>Before the expiration of the time limit for amending the claims and to be republished in the event of the receipt of amendments.</i>

(54) Title: ATM COMMUNICATIONS SYSTEM AND METHOD

MCA SSCS Packet Header (PH)

PPT - Payload Information Field
 CI - Change Indication
 SI - Sequence Indication
 RES - Reserved


(57) Abstract

There is disclosed a method and system for multiplexing data from a plurality of user data sources across an ATM adaptation layer type -2 connection, in which a multiplexed trunk group extends across a plurality of common part sub-layer protocol data unit (CPS-PDU) mini-cells, and across a plurality of ATM cells. Large trunk groups are assembled by use of a single bit continuation indicator in the service specific convergence sub-layer header (SSCS) of successive CPS-PDU mini-cells. A packet payload type field (PPT) of the common part sub-layer (CPS)/service specific convergence sub-layer (SSCS) is used to indicate timing of changes in number of user data sources in a trunk group and provides for robust error recovery on loss of a single CPS-PDU mini-cell.

FOR THE PURPOSES OF INFORMATION ONLY

Codes used to identify States party to the PCT on the front pages of pamphlets publishing international applications under the PCT.

AL	Albania	ES	Spain	LS	Lesotho	SI	Slovenia
AM	Armenia	FI	Finland	LT	Lithuania	SK	Slovakia
AT	Austria	FR	France	LU	Luxembourg	SN	Senegal
AU	Australia	GA	Gabon	LV	Latvia	SZ	Swaziland
AZ	Azerbaijan	GB	United Kingdom	MC	Monaco	TD	Chad
BA	Bosnia and Herzegovina	GE	Georgia	MD	Republic of Moldova	TG	Togo
BB	Barbados	GH	Ghana	MG	Madagascar	TJ	Tajikistan
BE	Belgium	GN	Guinea	MK	The former Yugoslav Republic of Macedonia	TM	Turkmenistan
BF	Burkina Faso	GR	Greece			TR	Turkey
BG	Bulgaria	HU	Hungary	ML	Mali	TT	Trinidad and Tobago
BJ	Benin	IE	Ireland	MN	Mongolia	UA	Ukraine
BR	Brazil	IL	Israel	MR	Mauritania	UG	Uganda
BY	Belarus	IS	Iceland	MW	Malawi	US	United States of America
CA	Canada	IT	Italy	MX	Mexico	UZ	Uzbekistan
CF	Central African Republic	JP	Japan	NE	Niger	VN	Viet Nam
CG	Congo	KE	Kenya	NL	Netherlands	YU	Yugoslavia
CH	Switzerland	KG	Kyrgyzstan	NO	Norway	ZW	Zimbabwe
CI	Côte d'Ivoire	KP	Democratic People's Republic of Korea	NZ	New Zealand		
CM	Cameroon		Republic of Korea	PL	Poland		
CN	China	KR		PT	Portugal		
CU	Cuba	KZ	Kazakhstan	RO	Romania		
CZ	Czech Republic	LC	Saint Lucia	RU	Russian Federation		
DE	Germany	LI	Liechtenstein	SD	Sudan		
DK	Denmark	LK	Sri Lanka	SE	Sweden		
EE	Estonia	LR	Liberia	SG	Singapore		

ATM COMMUNICATIONS SYSTEM AND METHOD
Field of the Invention

The invention relates to digital communications networks, and particularly, although not exclusively to an arrangement and method for transmitting multiplexed multi-user asynchronous transfer mode (ATM) traffic across such communications networks.

Background to the Invention

The known asynchronous transfer mode (ATM) transmission technique is a modern telecommunications switching technique which is able to switch connections for a wide range of different data types at a wide range of different bit rates. ATM technology provides a flexible form of transmission which allows various types of service traffic data, eg voice data, video data, or computer generated data to be multiplexed together onto a common physical means of transmission. Currently, several trends are encouraging the widespread introduction of ATM; for example the availability of high speed, low error rate communication links between switching centers, an availability of technology to digitize video and speech, and pressure to reduce operating costs by integrating previously separate telephony and data networks. ATM technology allows speech data, video data and inter-computer data to be carried across a single communications network. The information carried in each of these services is reduced to digitized strings of numbers which are transmitted across such a communications network from point to point.

Referring to Fig. 1 herein, there is shown schematically a portion of a communications network comprising first and second node devices 100, 101 respectively linked by a communications link 100. Transport of ATM data communications traffic is made across the communications link 102 between the first and second node devices, which may be for example switches 101, 102. Digitized data is received from customer equipment such as telephones, computers, faxes, modems and video broadcast apparatus in the form of frames of digitized signals at transmitting node, eg switch 100. The frames can either be of variable length or fixed length, and may arrive at the switch at a variable rate; or at a fixed rate. The frames of data arriving at the switch are packaged into ATM data cells 103, which have a fixed number of bytes. Transport

-2-

of ATM cells between node devices is handled by the node devices operating in accordance with the ATM protocol corresponding to the International Standards Organization (ISO) Open Systems Inter-connexion (OSI) architecture, layers two and three⁽¹⁾. Packaging of the incoming data frames received asynchronously from the customer equipment is handled by the switches operating in accordance with ATM adaptation layer (AAL) protocols which segment the arriving frames of data into payload data of the ATM cells at the transmission node, and reassemble the payload data into frames at the destination node 102.

10 The ATM adaptation layer corresponds to layer four of the OSI model. Equipment operating in accordance with the ATM adaptation layer protocols are capable of structuring incoming data in different ways, to suit different service types, eg video data, computer generated data, voice data. Many different service types can be implemented by the ATM adaptation layer simultaneously.

The basic reason for having ATM cells is that they have a fixed length. Fixed length cells are easier for hardware to handle than variable length frames. The ATM adaptation layer packages various types of data of variable length or fixed length frame type into the fixed length ATM cells for transport between physical devices. Because the ATM cell length was historically selected to accommodate various types of traffic, fixing the length of the ATM cell involved difficult decisions, and the final length of ATM cell selected is not perfect for each type of data carried.

25 The ATM cell comprises a header portion which carries routing information and other housekeeping information necessary for the operation of the ATM network, and a payload portion which carries the actual data traffic. To transfer delay sensitive services such as speech, it is important that the ATM cell be reasonably short in order to avoid unacceptably long delays in filling the cell payload portion before transmitting the cell across the network. On the other hand, for other types of traffic such as computer to computer file transfers longer cells are more efficient, since the proportion of available transmission bandwidth taken up by the ATM cell header compared to the data payload of the cell is reduced. For delay insensitive traffic, the overhead of the housekeeping information sent in the header of each ATM cell would be relatively large if short cells were to be used. Thus, the choice

-3-

of ATM cell size is a compromise and is settled at a length of 53 octets, comprising 48 octets of data payload (the ATM Service Data Unit, ATM-SDU) and a 5 octet header for transmission of housekeeping protocol information, as shown schematically in Fig. 2 herein. The protocol header in the ATM cell constitutes approximately 10% of the whole cell. This size of ATM cell introduces a delay in transmission of data which is significant for types of data having a low data rate, for example speech data. For example for a conventional 64 kilobits per second (kbits/s) voice data traffic, normal speech data samples are converted into one octet of digital data every 125 microseconds (μ s). Thus, $48 \times 125 \mu$ s = 6000 μ s are required to fill the 48 data octets of an ATM cell payload. This introduces a 6 millisecond (ms) delay to each cell transmitted, in addition to two network switching delays one from each switch, and transmission delays across the network. For speech services, it is important to have an effectively constant delay between source and destination of a call, and the delay must be reasonably short. Large variations in delay produce broken sound effects, and make voice signals unintelligible to a service user. Long delays, for example those sometimes encountered on transatlantic satellite links, make two-way conversation awkward. In general, a conventionally accepted maximum one-way delay for speech data is 25 ms. Delays longer than this, as well as making the speech service unacceptable to users, also require complicated and expensive echo suppression equipment, which has the additional disadvantage of introducing noise. Thus, the conventional 53 octet ATM cell is not ideal for 64 kbit/s voice data traffic. However, with the advent of mobile telecommunications systems, normal 64 kbit/s sampled voice signals are compressed using code compression algorithms, resulting in transmission data rates as low as 4 kbit/s. Under these circumstances, the delay introduced in filling a full ATM cell may be as high as 96 ms, an unacceptably high delay.

In view of the above delays and to accommodate different data traffic types, the ATM adaptation layer (AAL) is split into a number of sub-layers. A first sub-layer, the known AAL-type 1 layer is aimed at constant bit rate services. The currently developing, and not yet finalized AAL-type 2 layer (formerly known in Europe as AAL-type 6, and elsewhere as AAL-CU) allows multiple variable length sub-cells, called

mini-cells to be carried within one ATM cell. An object of AAL-type 2 is to support all services which require the multiplexing of information from multiple user data sources into a single ATM connection. The AAL-type 2 protocol, which breaks the basic rule of ATM that all cells be of fixed length, is aimed at being expedient for carrying low speed data where the delay caused by waiting for a full ATM cell to fill is too long, and the overhead of carrying an incomplete ATM cell is too great. However, the implementation of this layer is incomplete and there still remains a requirement for a method of transmitting data from a multiplicity of sources, including low bit rate data, over ATM networks in an efficient manner.

Summary of the Invention

One object of the present invention is to provide an improved method for transmitting delay sensitive data from a plurality of user data sources over a communications network.

Another object of the present invention is to provide a method of transmitting multi-service data from a range of different data sources over an ATM network in an efficient manner.

Another object of the present invention is to provide a robust method of transmitting large trunk groups of data of a plurality of data user sources, over a communications network.

According to one aspect of the present invention there is provided a method of transmitting data in a communications network, said method comprising the steps of:

receiving a plurality of data samples comprising at least one said data sample received from each of a plurality of user data sources;

multiplexing said plurality of data samples into a data payload of at least one asynchronous transfer mode mini-cell;

-5-

packetizing said mini-cell(s) into at least one asynchronous transfer mode cell; and

transmitting said asynchronous transfer mode cell(s).

5

Preferably each said mini-cell carries data from a plurality of user data sources.

According to a second aspect of the present invention there is provided a method of communicating user data of a plurality of user data sources over a communications network, said method comprising the steps of:

10

multiplexing a respective data sample from each of a plurality of user data sources to produce a frame of user data, said frame containing data of each of said plurality of user data sources; and

15

assembling said frame into a plurality of data payloads of a corresponding plurality of asynchronous transfer mode mini-cells.

20

Said frame may be of a length greater than a payload of a single said mini-cell.

A said frame may be partitioned to run consecutively across a plurality of said data packet payloads.

25

Preferably said method comprises the step of:

assembling a respective protocol header to each of said mini-cells, said protocol header comprising a continuation indicator signal indicating whether or not said frame continues beyond a length of said mini-cell.

30

Preferably said continuation indicator signal comprises a single bit field.

35

Preferably a said mini-cell comprises an asynchronous transfer mode adaptation layer-type 2 common part sub-layer packet.

Preferably said continuation indicator signal comprises an asynchronous transfer mode adaptation layer-type 2 header.

- 5 Preferably said continuation indicator signal comprises an asynchronous transfer mode adaptation layer- type 2 service specific convergence sub-layer field.

- 10 According to a third aspect of the present invention there is provided a method of communicating user data of a plurality of user data sources over a communications network, said method comprising the steps of:

15 signaling an asynchronous transfer mode adaptation layer type-2 connection;

 multiplexing user data of said plurality of user data sources into a trunk group frame;

- 20 assembling said trunk group frame into a payload of at least one asynchronous transfer mode variable length cell;

 transmitting said asynchronous transfer mode cell(s) across said network.

- 25 According to a fourth aspect of the present invention there is provided a method of communicating user data of a plurality of user data sources over a communications network, said method comprising the steps of:

30 multiplexing a data sample from each of a plurality of user data sources to produce a frame of user data, said frame containing data of each of said plurality of user data sources; and

- 35 assembling said frame into a data payload of at least one asynchronous transfer mode mini-cell.

Preferably said method comprises the step of:

including a protocol header signal in a said mini-cell, to indicate a change in number of user data sources who's data is assembled into a said at least one mini-cell.

5

Preferably said method comprises the step of using a packet payload type field of a service specific convergence sub-layer header of an asynchronous transfer mode mini-cell to indicate change of number of said user data sources who's data is carried in a series of said mini-cells.

10

Preferably said method comprises the step of using a packet payload type field of an asynchronous transfer mode mini-cell header to signal timing of a change of a number of said multiplexed users.

15

Preferably said mini-cell comprises an ATM adaptation layer type-2 mini-cell.

According to a sixth aspect of the present invention there is provided a method of communicating user data to a plurality of user data sources of a communications network, said method comprising the steps of;

20

multiplexing data of a first plurality of user data sources into a first data group having a first group size;

25

assembling said first data group into a data payload of a first set of at least one mini-cell;

30

multiplexing data of a second plurality of user data sources into a second data group having a second group size;

assembling said second data group into a data payload of a second set of at least one mini-cells,

35

wherein each said mini-cell comprises a respective protocol header, said protocol header arranged to indicate a change in group size between successive mini-cells.

According to a seventh aspect of the present invention there is provided a method of transmitting data in a communications network, comprising the steps of:

5

receiving a plurality of data samples comprising at least one said data sample received from each of a plurality of user data sources;

multiplexing said plurality of data sources into a trunk group;

10

establishing a trunk group connection using an asynchronous transfer mode adaptation layer-type 2 negotiation procedure;

15

signaling additions or subtractions of users in the trunk group by incorporation of signals contained within an asynchronous transfer mode adaptation layer-type 2 protocol header, whilst leaving said asynchronous transfer mode adaptation layer-type 2 trunk group connection intact.

20

According to an eighth aspect of the present invention there is provided a method of communicating user data of a plurality of user data sources across a communications network comprising a plurality of transmitting entities and receiving entities, said method comprising the steps of:

25

signaling to a said receiving entity to create a trunk group connection;

multiplexing data of said plurality of user data sources into a trunk group data frame;

30

signaling to said receiving entity a length of said trunk group data frame;

35

assembling said trunk group data frame into a plurality of asynchronous transfer mode adaptation layer type-2 mini-cells; and

signaling a change in length of said trunk group data frame in an asynchronous transfer mode adaptation layer type-2 mini-cell header.

5 Preferably said step of signaling a length of said trunk group data frame comprises signaling by an asynchronous transfer mode negotiation procedure protocol.

10 Preferably said step of signaling a change in length of said trunk group data frame comprises signaling within an asynchronous transfer mode adaptation layer type-2 payload packet type field.

15 According to a ninth aspect of the present invention there is provided a method of implementing changes in data payload of an asynchronous transfer mode adaptation layer type-2 mini-cell by decoding a payload packet type field of said mini-cell in accordance with the following steps:

20 decoding a received packet payload type signal 10 as indicating a mini-cell contains a new data payload containing data from a different number of user data sources compared with a previously received mini-cell.

25 Preferably said method comprises the step of decoding a received packet payload type field signals as follows:

30 decoding a first mini-cell having packet payload type field 00 followed by an immediately succeeding second mini-cell including packet payload type signal 00 as indicating no change in a number of user data sources in a data payload of said mini-cells.

35 Preferably said method comprises the step of decoding a received packet payload type field signals as follows:

decoding a first mini-cell having packet payload type field 01 followed by an immediately succeeding second mini-cell including packet payload type signal 01 as indicating no change in a number of user data sources in a data payload of said mini-cells.

Preferably said method comprises the steps of:

5 in an asynchronous transfer mode adaptation layer-type 2 connection, decoding a received packet payload type signal 00 of a first mini-cell and a received packet payload type signal 01 of a next received mini-cell as indicating a loss of data.

10 Preferably said method comprises the steps of:

 in an asynchronous transfer mode adaptation layer-type 2 connection, decoding a received packet payload type signal 01 of a first mini-cell and a packet payload type signal 00 of next received mini-cell
15 as indicating a loss of data.

Brief Description of the Drawings

20 For a better understanding of the invention and to show how the same may be carried into effect, there will now be described by way of example only, specific embodiments, methods and processes according to the present invention with reference to the accompanying drawings in which:

25 Fig. 3 illustrates conceptually a communications network hardware comprising first and second node devices connected by a communications link device, the node devices acting as receiving and transmitting entities for transmission of communications data;

30 Fig. 4 illustrates a transmission series of asynchronous transfer mode cells, containing a plurality of asynchronous transfer mode adaptation layer type-2 mini-cell (common part sub-layer packets);

35 Fig. 5 illustrates schematically a layout of an asynchronous transfer mode common part sub-layer packet mini-cell;

Fig. 6 illustrates schematically an asynchronous transfer mode mini-cell common part sub-layer protocol header, and a mini-cell data payload;

Fig. 7 illustrates schematically a start field header in an asynchronous transfer mode mini-cell data payload;

Fig. 8 illustrates schematically assembly of data of a plurality of user data sources into a plurality of CPS packet mini-cells in accordance with a single channel adaptation (SCA) method;

10

Fig. 9 illustrates schematically assembly of data from a plurality of user data sources into a plurality of CPS packet mini-cells in accordance with a multiple channel adaptation method (MCA) according to a specific method of the present invention;

15

Fig. 10 illustrates schematically a method of assembling a CPS packet mini-cell including a data payload comprising a multiplex of data of a plurality of user data sources, and assembly of a common part sub-layer/service specific convergence sub-layer (CPS/SSCS) header in accordance with a multiple channel adaptation method;

20

Fig. 11 illustrates schematically a plurality of CPS packet mini-cells containing a data payload of a single trunk group data frame comprising multiplexed data from a plurality of user data sources;

25

Fig. 12 illustrates schematically a series of CPS packet mini-cells transmitted across a communications network in accordance with specific methods of the present invention, and a result of decoding signaling information comprising the transmission, in accordance with further specific methods according to the present invention;

30

Fig. 13 illustrates schematically in overview, a signaling process according to a specific method of the present invention;

Figs. 14 to 16 illustrate schematically in overview, a signal decoding process comprising specific methods of the present invention;

35

Figs. 17 illustrates schematically a decoding sub-process according to a specific method of the present invention;

Fig. 18 illustrates schematically another decoding sub-process
5 according to a specific method of the present invention;

Fig. 19 illustrates schematically another decoding sub-process according to a specific method of the present invention;

10 Fig. 20 illustrates schematically a method of multiplexing long data structures into a plurality of mini-cells; and

Fig. 21 illustrates a method of multiplexing relatively short data structures into a plurality of mini-cells.

15

Detailed Description of the Best Mode for Carrying Out the Invention

There will now be described by way of example the best mode contemplated by the inventors for carrying out the invention. In the
20 following description numerous specific details are set forth in order to provide a thorough understanding of the present invention. It will be apparent however, to one skilled in the art, that the present invention may be practiced without using these specific details. In other instances, well known methods and structures have not been described in detail so as
25 not to unnecessarily obscure the present invention.

Referring to Fig. 3 herein, there are shown first and second ATM node device, 300, 301 respectively between which are transmitted data signals in the form of ATM cells 302 which are transmitted across a
30 physical transmission medium 303 between the node devices single user or multi-service data from a plurality of single service or multi-service data sources, for example speech data 304, video data 305, and computer generated data 306 are input into the node devices 300, 301 and incorporated into the plurality of ATM cells 302. Fig. 3 illustrates a simple
35 form of physical network comprising two node devices and one link device, in order to illustrate principles of the best mode for carrying out the invention without unnecessarily obscuring the invention. It will be

-13-

appreciated by the skilled reader that methods and processes according to the invention will be applicable to highly complex communications networks.

5 Referring to Fig. 4 herein, there is shown schematically a series of two conventional ATM cells, each containing a plurality of conventional ATM, CPS Protocol Data Unit (CPS-PDU) mini-cells MC1-MC6. A single mini-cell 400 comprises a 3 byte mini-cell header 401, and a variable length mini-cell payload 402 for carrying data traffic. In AAL-2 one or
10 more mini-cells may comprise the payload of an ATM cell, occupying a maximum of 48 octets of the ATM cell, the remaining 5 octets of the ATM cell being reserved for the ATM cell header. An ATM mini-cell may cross an ATM cell boundary. For example in Fig. 4 first ATM cell, cell 1 contains mini-cells MC1, MC2 and a portion of a further mini cell MC3.
15 The remaining portion of the mini-cell MC3 occupies a payload of second ATM cell, cell 2.

In AAL-type 2 the mini-cell header 401 is split into two parts; firstly a common part sub-layer (CPS) packet header and secondly a service
20 specific convergence sub-layer (SSCS) packet header. The common part sub-layer includes basic control information which is common to all mini-cells transmitted, such as a user identification (UID) indicating the number of users of a virtual channel, (an ATM virtual channel can have up to 256 individual users); a length indicator (LI) field indicating the
25 length of the mini-cell payload, the mini-cells being capable of having variable length payload; and a cyclic redundancy data (CRC) field to protect the content of the mini-cell header against errors. The mini-cell header is of length 3 bytes. Within the mini-cell header there are left five free bits allocated for the service specific convergence sub-layer (SSCS)
30 packet header information. The service specific convergence sub-layer packet header constitutes a user definable field, such that users can adapt the service specific convergence sub-layer header to tailor the ATM mini-cell to their own requirements. Users have an option to place service specific convergence sub-layer information as part of the mini-cell
35 payload 400, or can place the SSCS information in the vacant 5 bits in the mini-cell header 401. Thus, for example if a user wanted to include digital error protection information for a service, this could be included in

-14-

the payload of the mini-cell. In Figs. 5,6 and 7 herein, there is shown schematically the 3 byte mini-cell header, comprising the common part sub-layer (CPS), and the five free bits reserved for the service specific convergence sub-layer (SSCS), and also positioning of elements of the service specific convergence sub-layer at a location 501 in the mini-cell payload 400.

In the AAL-type 2, ATM cells are filled with CPS-PDU mini-cells, such that data from different users fill respective different mini-cells one user per mini-cell. Thus, using a single channel adaptation (SCA) method for example the from three different users can be transmitted in a single ATM cell, by incorporating three separate mini-cells as the payload of the ATM cell, each mini-cell containing data from a respective user data source. Thus, the ATM mini-cells can be transmitted asynchronously within the ATM cell structure, and a number of users may be multiplexed into the payload of an ATM cell, using a plurality of mini-cells. However, filling the mini-cells with data still incurs a packetization delay, as does filling an ATM cell. For transmission of very low bit rate services, which are delay sensitive, it is only feasible to include a small number of samples of the low bit rate data in the mini-cell payload, before the packetization delay becomes excessive. For example, at very low bit rates, it may only be possible to include 3 or 4 octet samples in a mini-cell payload before the packetization delay in filling the mini-cell becomes too great. The CPS packet header needs to be added to this payload, resulting in a mini-cell comprising a 3 byte CPS packet header, and 3 or 4 bytes of payload data. Thus, for low bit rate data, a 3 byte mini-cell header constitutes an excessive overhead for a 3 or 4 byte data payload, resulting in poor bandwidth utilization at low data rates.

Referring to Fig. 8 herein, there is illustrated packetization of multi-user data from a plurality of low bit rate synchronous user data sources using a SSCS-single channel adaptation mechanism. In SSCS-SCA, each mini-cell carries data of a single user data source. For example in Fig. 8, mini-cell 800 carries data from User 1 and mini-cell 824 carries data from User 24. Six bytes of traffic data 825-830 are encapsulated in a service specific convergence sub-layer packet header 831 to form a first service specific convergence sub-layer packet data unit SSCS-

-15-

PDU1. The SSCS-PDU1 comprising 6 bytes of user data from User 1, plus the SSCS packet header 831 are encapsulated in the CPS packet header 832 to form CPS packet mini-cell 801. In this case, the mini-cell 801 carries 6 bytes of user information with a minimum 3 byte packet header overhead. Similarly, for further users, User 2, to User 24 whose respective data are encapsulated in further respective mini-cells. The resulting CPS packet mini-cells are multiplexed together to form ATM cell payloads which are subsequently transmitted over the physical transmission media, the node devices 300, 301, and link 303 of the network operating in accordance with the ATM layer protocols. Using SSCS-SCA, information from each single user data source is used to construct a corresponding respective SSCS-SDU, containing that single user's data. To each SSCS-SDU is prepended a respective SSCS packet header, to form a series of SSCS-PDU's. To each of these is prepended a CPS packet header to form a respective set of CPS packet mini-cells. Each CPS packet mini-cell contains information from a single user, and so a unique ATM adaptation layer-type 2 connection is established for the support of each user. Thus, in the CPS packet header, the user ID (UID) is therefore associated with one particular SSCS user. This means that every time an AAL-type 2 connection is set up for a new user, an AAL-type 2 negotiation procedure (ANP) is invoked to negotiate a user ID (UID) value for assignment to that connection. Similarly, every time a user is released an AAL-type 2 negotiation procedure (ANP) is invoked to return that connection's user ID value to an unassigned state.

Whilst single channel adaptation is efficient for many service types, for some service types, in particular for low data rate services, the bandwidth utilization using SCA is inefficient.

30

Referring to Fig. 9 herein and hereafter generally, there is illustrated schematically another method of multiplexing a plurality of user data sources for transmission point to point across an ATM network in accordance with the best mode for carrying out the present invention. The specific methods and processes described hereafter not to be taken as limiting to the general scope of the invention. In Fig. 9, a plurality of network users, (User 1 to User 24), each generate synchronous user

35

data which is used to construct a multiple channel adaptation (MCA) CPS packet mini-cell in which the payload of the mini-cell represents a multiplex of the data from the multiple users. As shown in Fig. 9, a data octet 901 from a data source of User 1 is multiplexed with data octets from respective users; User 2, User 3 ... User 24, to form a payload of a first SSCS protocol data unit 902. Subsequent SSCS-packet data units carrying subsequent data octets of the plurality of uses are constructed similarly. For example where each user generates 6 bytes of data, for each user one byte of data is included in separate SSCS-protocol data units, SSCS-PDU's 1-6, culminating in the sixth SSCS-protocol data unit 907. In this manner, since a byte of data from each of the plurality of user data sources are multiplexed into a single SSCS-protocol data unit, before a next byte of information from the first user is considered, the first byte of information from each of the plurality of data sources can be packetized without waiting for the next byte of data to arrive from each or any user data source. The resultant SSCS protocol data units are prepended with a CPS packet header 908, resulting in CPS packets (mini-cells). These are multiplexed together to form the ATM cell payloads which are transmitted over the physical network. Each mini-cell is incorporated into one or more ATM cells for transmission across the network.

Thus, for example in the case of 24 users, using a multiple channel adaptation mini-cell (CPS-PDU-MCA), the mini-cell 902 comprises CPS packet header 908 and mini-cell payload (CPS packet payload) 909. The payload, rather than representing data from a single user, now represents multiplexed data from multiple users, so the user data is multiplexed on a frame by frame basis. Thus, if it is required to transmit user data of 24 different users from one point to another point, with each user generating an octet of data every 125 μ s, which would be the case for example for user data generated using the known pulse code modulation (PCM) method, then the 24 individual octet samples from the 24 users can be multiplexed into the payload of a single mini-cell. Information on the order in which the user data is packed into the mini-cell is not carried in the mini-cell header, but is transmitted separately via the ATM adaptation layer type 2 negotiation procedures (ANP). ANP messages are sent between two ATM adaptation layer type 2 entities (the node devices) to control assignment, removal, and status of AAL-

-17-

type 2 channels. Because the order in which the user data is packed into the mini-cell payload is known at the transmission end, and transmitted to the receiver entity in advance of transmitting the mini-cells, at the receiver end, the data can be unpacked from the mini-cell payload in known order and de-multiplexed into the 24 separate sets of user data.

Typical ANP messages include:

- Assignment request
- 10 • Assignment confirm
- Assignment denied
- Removal request
- Removal confirm
- Status poll
- 15 • Status response

For example, in Fig. 3 an assignment request message is sent by transmitting entity node device 300, and requests assignment of an AAL-type 2 channel. An assignment confirm message is sent by receiving entity, node device 301, and confirms assignment of an AAL-type 2 channel. An assignment denied message is sent by an AAL-type 2 entity which denies assignment of an AAL-type 2 channel. A removal request message is sent by an AAL-type 2 entity which requests removal of an AAL-type 2 channel. A removal confirm message is sent by an AAL-type 2 entity which confirms removal of an AAL-type 2 channel. A status poll message is sent by an AAL-type 2 entity which polls a status of an AAL-type 2 channel. A status response message is sent by an AAL-type 2 entity which responds to a status poll of an AAL-type 2 channel. Specific proposals for implementing the ANP protocol can be found in reference 2 herein. Information concerning the addition or deletion of users data sources in the multiplexed multi-user data package is signaled to the receiving entity using the ANP protocol as described, before the change of number of users actually occurs.

35 Referring to Fig. 10 herein, there is shown a way of structuring of an SSCS packet header within a CPS packet of an ATM mini-cell 1001 implementing a specific method according to the present invention. A

plurality N octets of data from a plurality N user data sources are multiplexed into data payload, service data unit packet 1002. SSCS packet header 1003 prepended to the service data unit 1002 comprises a 2 bit packet payload type field (PPT); a single bit change indication field (CI) 1005, a 2 bit Modulo 8 sequence indicator field (SN) 1006; and a 3 bit reserved field 1007. Using the SSCS packet header 1003 of Fig. 10 to prepend a multi-user multiplexed data packet 1002, information from multiple low bit rate sources can be multiplexed together into a trunk group data frame to achieve an improved bandwidth efficiency for a given bounded cell assembly delay. Irrespective of trunk group data frame size, the SSCS multi-channel access mechanism of Fig. 10 can be tuned to generate service data units of length approximately equal to the ATM cell payload size (maximum 48 bytes). This may yield an optimum balance between high bandwidth efficiency and elimination of potential error extensions. The size of the trunk group is deterministic, and at all times, a receiving entity node device has implicit knowledge of the length of the multi-user multiplexed packets which it receives. However, it is possible to change the size of the trunk group dynamically during the lifetime of a connection to accommodate changes in the number of low bit rate users. Changes in multi-user multiplexed data package size (ie the trunk group data frame) are signaled ahead using the ANP negotiation protocol, and the actual moment of implementation of the change of number of users is activated by transmission of the change indication (CI) bit signal within the SSCS packet header in which the change occurs, such that the receiving entity knows from the ANP negotiation process what the change in the number of users will be, and from the received mini-cells, the timing of when the change in number of users occurs.

Operation of the sequence indication (SI) field, continuation indication (CI) field and packet payload type (PPT) field will be described.

The 2 bit sequence indication field (SI) shown in Fig. 10 herein, is used in conjunction with the 1 bit start field sequence number (SN) in the start field of every CPS-packet data unit, as shown in Fig. 7 herein. The 2 bit sequence indication field (SI) provides a mechanism for guarding against the loss or mis-delivery of a mini-cell. For successive mini-cells

-19-

transmitted, the sequence indication field is incremented through the cycle 00, 01, 10, 11 and then back to 00. The receiving node device is configured to read the sequence indication field and check the cyclic incrementation of the sequence indication field. Any changes to the cycle
5 indicate that a mini-cell has been lost. Particularly in synchronous services, it is important to ensure that an end to end phase relationship is maintained, and detection of lost or miss-delivered mini-cells is important. Generally, loss of phase may lead to error extension, particularly for modem traffic, where the modem requires a significant duration in order
10 to regain synchronization, once synchronization is lost.

The single bit continuation indication field (CI) can be used to incorporate a large trunk group data frame, ie a long sequence of multiplexed users, into a plurality of mini-cells. For example, in Fig. 11
15 herein, 160 octets from 160 user data sources respectively are assembled into a single trunk group data frame. Referring to Fig. 11 herein, first to fourth successive CPS packet mini-cells 1101-1104 respectively each have a respective data payload which is filled with the data octets from the multiple users. The trunk group is transmitted in the
20 first to fourth mini-cells 1101-1104 by filling data payload of the first mini-cell 1101 with octets from users 1 to 45, filling data payload of the second mini-cell 1102 with octets from data users 46 to 90, filling data payload of third mini-cell 1103 with octets from data users 91 to 134, and filling data payload of the fourth mini-cell 1104 with octets from data users 136 to
25 160. In the SSCS packet header 1003 of each mini-cell, the single bit continuation indication field (CI) can adopt a value of 0 or 1. The value 1 is used to signify that at the end of the mini-cell containing the continuation indicator 1, the data payload continues into a subsequent mini-cell. Thus, in Fig. 11 first mini-cell 1101 containing continuation
30 indicator value 1 has a data payload part of the trunk group data frame which continues into subsequent second mini-cell 1102. Similarly for mini-cells 1102 and 1103, which also contain continuation indicator value 1. The fourth mini-cell contains a continuation indicator of value 0, indicating that the data payload of the fourth continuation cell does not
35 continue into a subsequent fifth mini-cell.

-20-

For mini-cells shown in Fig. 11 a maximum data payload of 45 octets is selected as a default condition. This value has been selected to coincide with the suggested default length of mini-cell stated in reference 2 herein. If very large trunk group payloads are assembled in a single large mini-cell, then several successive ATM cell payloads could be generated without any CPS protocol control information. Under these circumstances, a single error in the data payload could lead to a prolonged error extension, thus degrading the overall error performance of the ATM communications link. In the present method, although the maximum mini-cell payload size of 45 octets is complied with, there is no restraint placed on the maximum size of a trunk group itself which can be transmitted. Mini-cells having a continued payload, and having a continuation indicator value of 1 are set at single constant length of the default mini-cell payload size of 45 octets. For mini-cells having payloads which are not continued to a successive mini-cell, the size of the mini-cell payload is variable to accommodate the remainder of the trunk group data. Since the receiving entity has an implicit knowledge of the length of the trunk group, due to prior signaling via the ANP mechanism, the lengths of both the continued mini-cell payloads, and the final un-continued mini-cell payload are known implicitly. However, whilst the best mode herein contemplates a mini-cell of maximum payload 45 octets, the invention is not restricted to such mini-cell payload length. In a further specific method according to the present invention, the mini-cell payload size of a continued mini-cell may be set at an optional maximum mini-cell payload length of 64 octets.

The two bit packet payload type field (PPT) is used both to designate the type of packet payload, whether operation and maintenance (OAM) data, or user data, and to dynamically indicate in which mini-cell a change of number of users in the trunk group occurs. In the 2 bit packet payload type field, the value 11 is assigned as indicating that the mini-cell contains operation and maintenance signal information (OAM). In this case, the data payload of the mini-cell contains operation and maintenance signal information, and not trunk group user data. The other three states of the 2 bit packet payload type (PPT) field are allocated as follows:

-21-

- **00** - indicates an SSCS packet which carries user information (ie trunk group information of the multiplexed multi-user data).
- 5 **01** - indicates an SSCS packet carrying user information, the same as value 00 above.
- Values **00** and **01** are used to alternate with each other each time there is a change in number of users carried over a connection.
- 10
- **10** indicates an SSCS packet which is the first packet of a group of SSCS packets containing data of a different number of users as the immediately preceding SSCS packet. In otherwords **10** indicates a first SSCS packet containing data of a new plurality of users.
- 15

Referring now to table 1 herein, there is shown a state table which may be used to decode the packet payload type (PPT) fields when a change in trunk group structure size has occurred. In the table "same" and "inverted" imply that the PPT flag is equal to **00** or **01**. Same means that it is the same value as the last received user mini-cell, whilst inverted means it is the opposite value.

25 Table 1

P1	SN	Action
Same	X+1	Normal operation - no structure size change detected
Same	X+2	Missing mini-cell detected - no structure size change detected
10	X+1	Structure size change detected - new size start this mini-cell
10	X+2	Missing mini-cell detected - structure size change detected, started new mini-cell

Inverted	X+2	Missing mini-cell detected - structure size change detected, started in previous mini-cell
All other combinations		Errored condition - take mitigating action as necessary, ie verify/resync at next mini-cell

Since the values **00** and **01** are both used to indicate SSCS packets which carry user information, but are two different states, a transition from **00** to **01** or from **01** to **00** is used to indicate a permanent change in the size of the mini-cell payload has occurred. For example. Referring to Fig. 12 herein where mini-cells A to H are transmitted, and the number of users in a trunk group varies from 20 users, to 21 users and then to 22 users, the packet payload type indicator values **00**, **01** are used to indicate permanent changes of user number as follows. Firstly the change in user number from 20 users to 21 users is signaled from the transmitter node entity to the receiver node entity in advance of the change, using the ATM adaptation layer type-2 negotiation procedure (ANP). The number of users in the trunk group is then changed at the transmitter end, and the new number 21, of users, are included in the trunk group. In the first mini-cell (mini-cell D) carrying data from the new trunk group of 21 users, the user data payload of the mini-cell is packetized by adding the SSCS header, including in the packet payload type field change pulse signal, **10**. At the receiver, a change pulse decoder apparatus decodes the received SSCS packet header and determines from the PPT value **10** that a change in user number has occurred in the trunk group, starting at mini-cell D. The receiver has information in advance of mini-cell D, of the actual number of users in the new trunk group, since this has been transmitted earlier using the ANP protocol. On receiving the change pulse signal PPT value **10** the receiver may then allocate the 21 users octets to a respective 21 user end point destinations, taking the timing of this allocation from the change pulse indicator value **10** of the packet payload type field.

In the next mini-cell, mini-cell E the packet payload type field (PPT) changes to the alternate value of the permanent change indicator (PCI), ie **01**. The permanent change indicator PPT value **01** indicates that a change in the number of users of the trunk group has been made. The

-23-

value of packet payload type is maintained at the new value of permanent change indicator PPT value, **01** for as long as the number of users in the trunk group remains the same, in this case up to mini-cell F.

5 New data concerning a new number of users of the trunk group is signaled ahead using the ANP protocol to the receiver, during transmission of mini-cells D to F, to enable the receiver to set up for a new change of trunk group user at mini-cell G. In this example, mini-cells G, H include trunk data from 22 users, effective as from mini-cell G. Data
10 from the new number of user data sources of the trunk group are multiplexed into a trunk group frame which is transmitted in mini-cells G, H and subsequent mini-cells. At the beginning of the first mini-cell containing the amended number of trunk group users, ie mini-cell G, the PPT field of mini-cell G contains the value **10**, being the change pulse CP
15 signal. At the receiver, the PPT field value **10** is decoded as indicating that the new number of users in the trunk group is effective as from mini-cell G. The decoder proceeds to de-multiplex mini-cell G and subsequent mini-cells in accordance with the revised ANP information indicating that there are now 22 users in the trunk group and received
20 multiplexed data octets of the 22 user data sources are sent to the corresponding respective 22 user destinations at the receiver switch. In the next mini-cell following the mini-cell G containing the change pulse, the PPT field reverts to the other permanent change indicator value **00**.

25 Referring to Fig. 13 herein, a general overview of a transmission method is shown. In step 1301, the transmitter signals ahead to the receiver, indicating a new trunk group size, using the ANP protocol. In step 1302, the number of users of the trunk group are changed and data from the new plurality of users are multiplexed into a new trunk group
30 frame. In step 1303, the trunk group frame containing the new plurality of user data is packetized into one or more mini-cell payloads. To each mini-cell payload is added an SSCS field header. To the first mini-cell which carries the first trunk group frame having a new plurality of users, there is added in the PPT field the change pulse, value **10** in step 1304.
35 In step 1305 to subsequent mini-cells containing information from the same trunk group frame, there is added the permanent change indicator value of **00** or **01**. The value **00** or **01** is selected as being a different

value to the previous permanent change indicator value used for mini-cells carrying precious trunk group frames of the previous, different plurality of users.

5 Referring to Fig. 14 herein, there is shown a general overview of a process implemented at the receiver entity for decoding the mini-cell header information and for allocating mini-cell data payloads to end user destinations. In step 1401, the receiving entity receives a sequence of mini-cells. In step 1402, the receiver decodes the PPT field in the SSCS
10 header of the incoming mini-cells. If the decoder detects no change in the PPT field compared to a previously received PPT field of the previously received mini-cell, then in step 1404, the decoder demultiplexes the data payload of the mini-cell and switches the data octets of that data payload of the mini-cell to the plurality of user end
15 destinations specified in the currently held ANP information at the receiving entity. However, if in step 1403, the decoding receiver detects a change in the PPT field compared to the PPT field of the previously received mini-cell, then in steps 1405-1408, the receiving switch ascertains the value of the PPT field, either **00**, **01**, **10** or **11**.

20

If the PPT field value is **11**, then the receiving switch treats the data payload of the mini-cell containing the PPT field value as operation and maintenance (OAM) data in step 1409.

25 If the receiving switch decodes the PPT value of its currently received mini-cell as being **10** in step 1407, then the switch implements procedure upon change of number of users in a trunk group as illustrated in Fig. 15 herein. If the receiving switch decodes the PPT field of its current received mini-cell as being value **00** or value **01** in steps 1405 or
30 1406, then the receiving switch checks whether the previous PPT value of the previous mini-cell was **10**, indicating a change in number of users in the trunk group frame. If the PPT value in the previous mini-cell was **10** in step 1411, then the receiving switch treats the present mini-cell as the previous mini-cell, and de-multiplexes its data payload in accordance
35 with the current information received from the ANP protocol, and distributes the data octets of the data payload of the current mini-cell to the appropriate plurality of users indicated in the ANP protocol

-25-

information in step 1404. However, if the current mini-cell has a PPT value of **00** or **01**, and the PPT value of the previous mini-cell was not **10**, since in step 1403 it has been checked whether the PPT field has changed from that contained in the previous mini-cell, this indicates that
5 an error has occurred, and one or more mini-cells have been lost during transmission. A lost mini-cell recovery procedure 1412 illustrated in Fig. 16 is then followed.

Referring to Fig. 15 herein, where the PPT value **10** is received in
10 the current mini-cell at the receiving entity, this indicates that a change in number of users in the trunk group has occurred starting in the currently received mini-cell in step 1501. The receiving switch checks whether the latest ANP information has already been implemented yet in step 1502. If the latest ANP information has not yet been implemented, then the PPT
15 value **10** indicates the timing of the change to the new plurality of users indicated by the latest received ANP information and in step 1503 the receiving switch de-multiplexes the current mini-cell data payload and sends data octets of the new plurality of users indicated in the ANP information to a corresponding respective new number of user
20 destinations. If the latest ANP information received by the switch has been implemented in step 1502, then an error has occurred.

Referring to Fig. 16 herein, where the PPT values **00** or **01** are received, and this is a change from a PPT value in the previous mini-cell,
25 and the value of the PPT field in the previous mini-cell was not **10**, then this indicates that one or more mini-cells have been lost in transmission. In general, where mini-cells are lost during transmission this can cause de-synchronization of certain types of apparatus, for example modems, which take a large number of cycles to re-synchronize. Loss of a single
30 mini-cell due to cell congestion can occur relatively commonly compared to other cell loss mechanisms. Thus, recovery of single mini-cell loss errors may improve the performance of customer equipment attached to ATM networks, by avoiding loss of synchronization. Where a large number of mini-cells are lost, for example more than one it may be that
35 synchronization of modem equipment is lost in any case. In the best mode described herein, the PPT field can be used to recover single mini-cell loss errors as described in Fig. 16 herein. In step 1601, the 2 bit

-28-

cell loss conditions, and thus meets at least the following minimum requirements:

- 5 1) An ability to determine a phase/start of structure size change.
- 2) An ability to always determine current structure size even in the event of burst error conditions, ie it is always possible to attain full re-synchronization even when a change is pending.
- 10 3) No error extension in the event that the change occurs in the presence of a single cell loss/error condition.

Requirement 2) above dictates that the basic mechanism should provide a permanent indication that the change has occurred. Any
15 structure change relying solely on a transitory indicator could be missed completely under cell loss conditions, making the resultant re-synchronization process complex and potentially ambiguous.

Requirement 3) above is more demanding. In the event of a packet
20 loss a permanent change indicator will generally indicate that a change is about to occur or has already occurred. It may not be sufficient to predict the exact phase boundary of the change. If this is misinterpolated, in the worst case a permanent phase change will occur. To meet this requirements the in-band structure size change indication is implemented
25 via codes with the PPT as described herein, which provide both a permanent and a transitory indication of the change.

The **00** and **01** values of the PPT are used to provide the permanent change indication - the value of the PPT is flipped between
30 these values between successive changes of the trunk group structure size. The transitory element is provided using the **10** value of PPT field, giving further correlation of the structure size change. The PPT value is "pulsed" once to indicate the position of the change in the trunk group frame. For example if the original value of the PPT is **00**, then all user
35 mini-cells will contain this value up to the mini-cell containing the start of the structure size change. This mini-cell alone will contain a PPT value of **10**, thereafter (until the next change in trunk group size is made) the mini-

cells will contain a PPT set to 01. The mechanism may provide a secure change indication method even in the event of cell loss.

As can be seen from the foregoing, a difference between the single
5 channel adaptation mechanism and the multiple channel adaptation
mechanism as described herein lies in the functionality of the SSCS. For
AAL-2 SCA, the SSCS uses information from a single source to generate
each SSCS-packet data unit whereas for AAL-type 2 multiple channel
adaptation, the SSCS uses information from multiple sources to
10 generate an SSCS-packet data unit. The single channel adaptation and
multiple channel adaptation packets may be transported together on the
same ATM connection. The AAL-2 multiple channel adaptation
mechanism resides within the SSCS sub-layer and may be implemented
at no cost to other applications. In addition, it places minimal extra
15 requirements on the AAL-2 negotiation procedures (ANP).

The low bit rate synchronous services supported by SCA and MCA
comprise any low bit rate (64 kbits/s or lower) service which generates
user information on a fixed periodic basis. For example, for 64 kbits/s
20 PCM 32 kbits/s ADPCM, and 16 kbits/s ADPCM, an octet of information
is generated every 125 μ S, 250 μ S, and 500 μ S respectively. For LD-
CELP 10 bits of information are generated every 625 μ S. Table 2 below
shows that for these services, by multiplexing on a trunk group basis
using SSCS-MCA a considerable increase in bandwidth utilization
25 efficiency for a given bounded cell assembly delay can be attained
compared with SSCS-SCA. The results in table 2 assume a 1 μ S SSCS-
PDU assembly delay.

Table 2

Coding Algorithm	Efficiency	
	SSCS-MCA	SSCS-SCA
64 kbits/s PCM	86-92%	71%
32 kbits/s ADPCM	86-92%	56%
16 kbits/s ADPCM	86-92%	39%
LD-CELP	91%	25%

There are a number of potential applications than can be readily identified that will benefit from the use of MCA . These include PBX to PBX trunking, MSC to MSC trunking, variable P_xP_x sub-rate for multi-media to the desk top, and Legacy inter-working to the public switch telephone network (PSTN). The ability to multiplex on a trunk group basis as provided by SSCS-MCA together with the ability to multiplex on the user by user basis using SSCS-SCA may significantly enhance the applicability and flexibility of the AAL-type 2 layer.

SSCS-MCA methods described herein may enable information from multiple low bit rate sources to be multiplexed together on a trunk group basis to achieve a high bandwidth efficiency for a given bounded cell assembly delay. Irrespective of trunk group size, the SSCS-MCA mechanism can be tuned to generate SDU's of length approximately equal to the ATM payload size. This may yield an optimum balance between high bandwidth efficiency and the elimination of potential error extension. At all times the receiving station may have implicit knowledge of the length of the packets which it receives. However, it is possible to change the size of the trunk group dynamically during the lifetime of a connection to accommodate changes in the community of low bit rate users. The structure size change is performed in a controlled manner through a combination of the ANP negotiation procedure and in-band synchronization, and so again the receiving entity has full knowledge of the new structure before the change is made.

The SSCS-MCA mechanism may be used to achieve optimum bandwidth utilization whilst minimizing error extension effects and enabling MCA users to be freely multiplexed with SCA users. This can

be achieved by tuning the MCA mechanism such that it produces SDU's of near optimum length irrespective of trunk group size. The optimum packet length is equal to the free packet payload size of the CPS-PDU thus minimizing CPS-packet header overhead whilst still guaranteeing one CPS packet header per cell to avoid error extension effects. This can be achieved in two ways:

For large trunk groups (of length greater than the CPS packet payload size) SSCS-MCA can segment trunk group frame the structure across multiple SDU's; and for small trunk groups it is possible to concatenate several frames of data into one SDU.

Using SSCS-MCA recovery (without loss of synchronization) may be attainable after the loss of a single mini-cell or other error condition. There is no requirement for error detection or correction over the payload information in the SSCA-MCA SDU. It is possible to determine when there is an error in the SSCS packet header. The error control field acting over the CPS packet header is sufficient to further minimize any risk of possible misconnection due to error in the UUI field.

20

Referring to Fig. 17 herein, there is shown an example of how a trunk group frame size change can be interpolated during the loss of a single mini-cell, whilst still maintaining synchronization. In Fig. 17, there is shown a case where a CPS packet mini-cell sequence incurs loss of a CPS packet before a change of trunk group size. The packet sequence is incremented using the sequence indicator 5, 6, 7, 0, 1, 2, 3, 4. The PPT field values during the sequence of 8 transmitted packets is 01, 01, 01, 10, 00, 00, 00, 00. The third packet is lost (packet sequenced 0), but because the packet, sequenced 1, has a PPT field set to 10 the receiver infers that the lost packet, sequenced 0, did not contain a structure size change, but the received packet sequenced 1 does contain a structure size change.

Referring to Fig. 18 herein, a sequence of eight mini-cells (CPS packets) is transmitted. In this case, a packet containing the change pulse (CP) is lost. The packets are sequenced using the sequence indicator and sequence number in a sequence 5, 6, 7, 0, 1, 2, 3, 4 and

the fourth CPS packet, sequenced 0, is lost during transmission. Since the next packet, packet 1 has an inverted permanent change indicator compared to the last received packet 7, received before the lost packet, the receiver can infer that a change must have occurred during the lost packet.

Referring to Fig. 19 herein, there is shown an error recovery procedure in a case where a CPS packet mini-cell is lost after a change of number of users in the trunk group. Mini-cells are sequenced 5, 6, 7, 0, 1, 2, 3, 4 and CPS packet mini-cell 1 is lost during transmission. Since mini-cell 1 is received, the trunk group user number size change is implemented by the receiving switch. Mini-cell 2 is lost, but mini-cell 3 has an inverted permanent change indicator (PCI value of 00 compared to the previous permanent change indicator) PCA value 01 before the change in trunk group size. The receiver can therefore infer that there have been no changes in trunk group size contained in lost packet 2.

Referring to Fig. 20 herein, there is illustrated an example of segmentation of a large trunk group containing a large plurality of user data from a large plurality of user data sources into several successive mini-cells. There is no restraint on the maximum size of trunk group which can be multiplexed into a number of cells in the present best mode, since the 1 bit continuation indication field (CI) can be used to assemble large trunk groups into plurality of mini-cells as hereinbefore described with reference to Fig. 11.

Fig. 20 illustrates multiplexing of data from 49 user data sources into a plurality of mini-cell packets. A user trunk group containing 49 octets of user data from 49 different user data sources is multiplexed into a plurality of mini-cells. The first mini-cell contains the continuation indicator (CI) value 1, indicating that the data payload of the mini-cell is continued into the next mini-cell 2002. The first mini-cell 2001 includes a PPT field value 01 (this value could have also been value 00, depending upon the previous inversion 00 or 01 in previous mini-cells). The first 45 octets of data are included as the data payload of first mini-cell 2001. The CPS packet header of the first mini-cell 2001 includes a length indicator of 45 (LI = 45). The trunk group data is continued into second

mini-cell 2002. The continuation indicator in the second mini-cell 2002 is set at 0, indicating that the data payload of the second mini-cell is not continued into a third mini-cell. The second mini-cell includes channels 46 to 49 of the user data, ie a relatively short data payload. The length indicator in the corresponding CPS packet header is set to 4, indicating a payload size of 4 octets. The next trunk group frame of 49 users is multiplexed in a similar way in third mini-cell 2003 and subsequent mini-cells.

Referring to Fig. 21 herein, there is shown an example of multiplexing of a relatively small trunk group having a relatively small number of user data sources. Several successive trunk group frames can be concatenated into a single CPS packet payload. Since packet assembly delay is increased when successive trunk group frames are multiplexed together, a minimum packet assembly delay is therefore controlled by specifying a minimum trunk group size. For example a minimum trunk group size of 6 users implies that it will take almost 7 successive frames to generate an SSCS-service data unit whose length will fit into a maximum length CPS packet payload. At a 64 kbit/s user bit rate, this implies a maximum packet assembly delay of less than 1 ms, which is satisfactory for many applications. The maximum packet assembly delay can be increased or decreased by defining a minimum trunk group size accordingly. There is no net penalty in delay with this approach. Concatenation of structures into an SSCS protocol data unit first is equivalent to multiplexing packets by the CPS into an ATM cell.

The number of frames that can be bound into a single SSCS data unit is dependent upon the current trunk group structure size. It may be calculated as the integer division of the maximum CPS packet payload size (45 octets) by the trunk group size. The CPS packet header overhead is thus minimized whilst the resulting protocol data unit does not exceed the maximum packet length limit. For example with a 6 channel trunk group this implies that 7 successive frames can be concatenated together. The efficiency of the SSCS-MCA connection therefore varies with a structure size, but in all cases the utilization is significantly higher than that attained for a similar SSCS single channel adaptation connection. Further, although the size of the packet varies

with the group structure size, its size is completely deterministic and the receiver always has implicit knowledge of its length.

Abbreviations

5

AAL	ATM adaptation layer
AAL-type 1	ATM adaptation layer type 1
AAL-type 2	ATM adaptation layer type 2
ANP	ATM negotiation procedure
ATM	asynchronous transfer mode
CI	continuation indication field (1 bit)
CP	change pulse
CPS	common part sub-layer
CRC	cyclic redundancy data
LI	length indication
MCA	multiple channel adaptation
OAM	operation and maintenance
PCI	permanent change indicator
PDU	protocol data unit
PPT	packet payload type
SCA	single channel adaptation
SDU	service data unit
SI	sequence indication field (2 bits)
SN	sequence number (1 bit) in start field
SSCS	service specific convergence sub-layer
UID	user identification
UUI	CPS user to user indication

References

- [1] Copies of the ATM standards protocols are available from International Telecommunications Union (ITU), Sales and Marketing Service, Place des Nations, CH-1211, Geneva 20, Switzerland, telephone +41 22 730-6666 or from the ATM Forum, 2570 West El Camino Real, Suite 304, Mountain View, California CA90404, USA.

10

-35-

[2] Draft ITU-T Recommendation I.363.2 "B-ISDN ATM Adaptation Layer Type-2 Specification" (Madrid 1996), Recommendation I.363.2 (November 1996), available from International Telecommunications Union.

Claims:

1. A method of transmitting data in a communications network, said method comprising the steps of:

5 receiving a plurality of data samples comprising at least one said data sample received from each of a plurality of user data sources;

10 multiplexing said plurality of data samples into a data payload of at least one asynchronous transfer mode mini-cell;

packetizing said mini-cell(s) into at least one asynchronous transfer mode cell; and

15 transmitting said asynchronous transfer mode cell(s).

2. A method as claimed in claim 1, wherein each said mini-cell carries data from a plurality of user data sources.

20 3. A method of communicating user data of a plurality of user data sources over a communications network, said method comprising the steps of:

25 multiplexing a respective data sample from each of a plurality of user data sources to produce a frame of user data, said frame containing data of each of said plurality of user data sources; and

assembling said frame into a plurality of data payloads of a corresponding plurality of asynchronous transfer mode mini-cells.

30 4. A method as claimed in claim 3, wherein said frame is of a length greater than a payload of a single said mini-cell.

35 5. A method as claimed in claim 3 or 4, wherein a said frame is partitioned to run consecutively across a plurality of said data packet payloads.

-37-

6. A method as claimed in any one of claims 3 to 5, comprising the step of:

5 assembling a respective protocol header to each of said mini-cells, said protocol header comprising a continuation indicator signal indicating whether or not said frame continues beyond a length of said mini-cell.

7. A method as claimed in claim 6, wherein said continuation indicator signal comprises a single bit field.

10

8. A method as claimed in any one of claims 3 to 7, wherein a said mini-cell comprises an asynchronous transfer mode adaptation layer-type 2 common part sub-layer packet.

15 9. A method as claimed in any one of claims 3 to 8, wherein said continuation indicator signal comprises an asynchronous transfer mode adaptation layer-type 2 header.

20 10. A method as claimed in claim 9, wherein said continuation indicator signal comprises an asynchronous transfer mode adaptation layer-type 2 service specific convergence sub-layer field.

25 11. A method of communicating user data of a plurality of user data sources over a communications network, said method comprising the steps of:

signaling an asynchronous transfer mode adaptation layer type-2 connection;

30 multiplexing user data of said plurality of user data sources into a trunk group frame;

assembling said trunk group frame into a payload of at least one asynchronous transfer mode variable length cell;

35

transmitting said asynchronous transfer mode cell(s) across said network.

-38-

12. A method of communicating user data of a plurality of user data sources over a communications network, said method comprising the steps of:

5

multiplexing a data sample from each of a plurality of user data sources to produce a frame of user data, said frame containing data of each of said plurality of user data sources; and

10 assembling said frame into a data payload of at least one asynchronous transfer mode mini-cell.

13. A method as claimed in claim 12, comprising the step of:

15 including a protocol header signal in a said mini-cell, to indicate a change in number of user data sources who's data is assembled into a said at least one mini-cell.

14. A method as claimed in claim 12 or 13, comprising the step
20 of using a packet payload type field of a service specific convergence sub-layer header of an asynchronous transfer mode mini-cell to indicate change of number of said user data sources who's data is carried in a series of said mini-cells.

25 15. A method as claimed in any one of claims 12 to 14, comprising using a packet payload type field of an asynchronous transfer mode mini-cell header to signal timing of a change of a number of said multiplexed users.

30 16. A method as claimed in any one of claims 12 to 15, wherein said mini-cell comprises an ATM adaptation layer type-2 mini-cell.

17. A method of communicating user data to a plurality of user data sources of a communications network, said method comprising the
35 steps of;

-39-

multiplexing data of a first plurality of user data sources into a first data group having a first group size;

assembling said first data group into a data payload of a first set of
5 at least one mini-cell;

multiplexing data of a second plurality of user data sources into a second data group having a second group size;

10 assembling said second data group into a data payload of a second set of at least one mini-cells,

wherein each said mini-cell comprises a respective protocol header, said protocol header arranged to indicate a change in group size
15 between successive mini-cells.

18. A method of transmitting data in a communications network, comprising the steps of:

20 receiving a plurality of data samples comprising at least one said data sample received from each of a plurality of user data sources;

multiplexing said plurality of data sources into a trunk group;

25 establishing a trunk group connection using an asynchronous transfer mode adaptation layer-type 2 negotiation procedure;

signaling additions or subtractions of users in the trunk group by incorporation of signals contained within an asynchronous transfer mode adaptation layer-type 2 protocol header, whilst leaving said asynchronous
30 transfer mode adaptation layer-type 2 trunk group connection intact.

19. A method of communicating user data of a plurality of user data sources across a communications network comprising a plurality of
35 transmitting entities and receiving entities, said method comprising the steps of:

-40-

signaling to a said receiving entity to create a trunk group connection;

5 multiplexing data of said plurality of user data sources into a trunk group data frame;

signaling to said receiving entity a length of said trunk group data frame;

10 assembling said trunk group data frame into a plurality of asynchronous transfer mode adaptation layer type-2 mini-cells; and

15 signaling a change in length of said trunk group data frame in an asynchronous transfer mode adaptation layer type-2 mini-cell header.

20. A method as claimed in claim 19, wherein said step of signaling a length of said trunk group data frame comprises signaling by an asynchronous transfer mode negotiation procedure protocol.

20 21. A method as claimed in claim 19 or 20, wherein said step of signaling a change in length of said trunk group data frame comprises signaling within an asynchronous transfer mode adaptation layer type-2 payload packet type field.

25 22. A method of implementing changes in data payload of an asynchronous transfer mode adaptation layer type-2 mini-cell by decoding a payload packet type field of said mini-cell in accordance with the following steps:

30 decoding a received packet payload type signal 10 as indicating a mini-cell contains a new data payload containing data from a different number of user data sources compared with a previously received mini-cell.

35 23. A method as claimed in claim 22, comprising the step of decoding a received packet payload type field signals as follows:

decoding a first mini-cell having packet payload type field 00 followed by an immediately succeeding second mini-cell including packet payload type signal 00 as indicating no change in a number of user data sources in a data payload of said mini-cells.

5

24. A method as claimed in claim 22, comprising the step of decoding a received packet payload type field signals as follows:

10 decoding a first mini-cell having packet payload type field 01 followed by an immediately succeeding second mini-cell including packet payload type signal 01 as indicating no change in a number of user data sources in a data payload of said mini-cells.

15

25. A method as claimed in claim 22, comprising the steps of:

in an asynchronous transfer mode adaptation layer-type 2 connection, decoding a received packet payload type signal 00 of a first mini-cell and a received packet payload type signal 01 of a next received mini-cell as indicating a loss of data.

20

26. A method as claimed in claim 22, comprising the steps of:

25 in an asynchronous transfer mode adaptation layer-type 2 connection, decoding a received packet payload type signal 01 of a first mini-cell and a packet payload type signal 00 of next received mini-cell as indicating a loss of data.

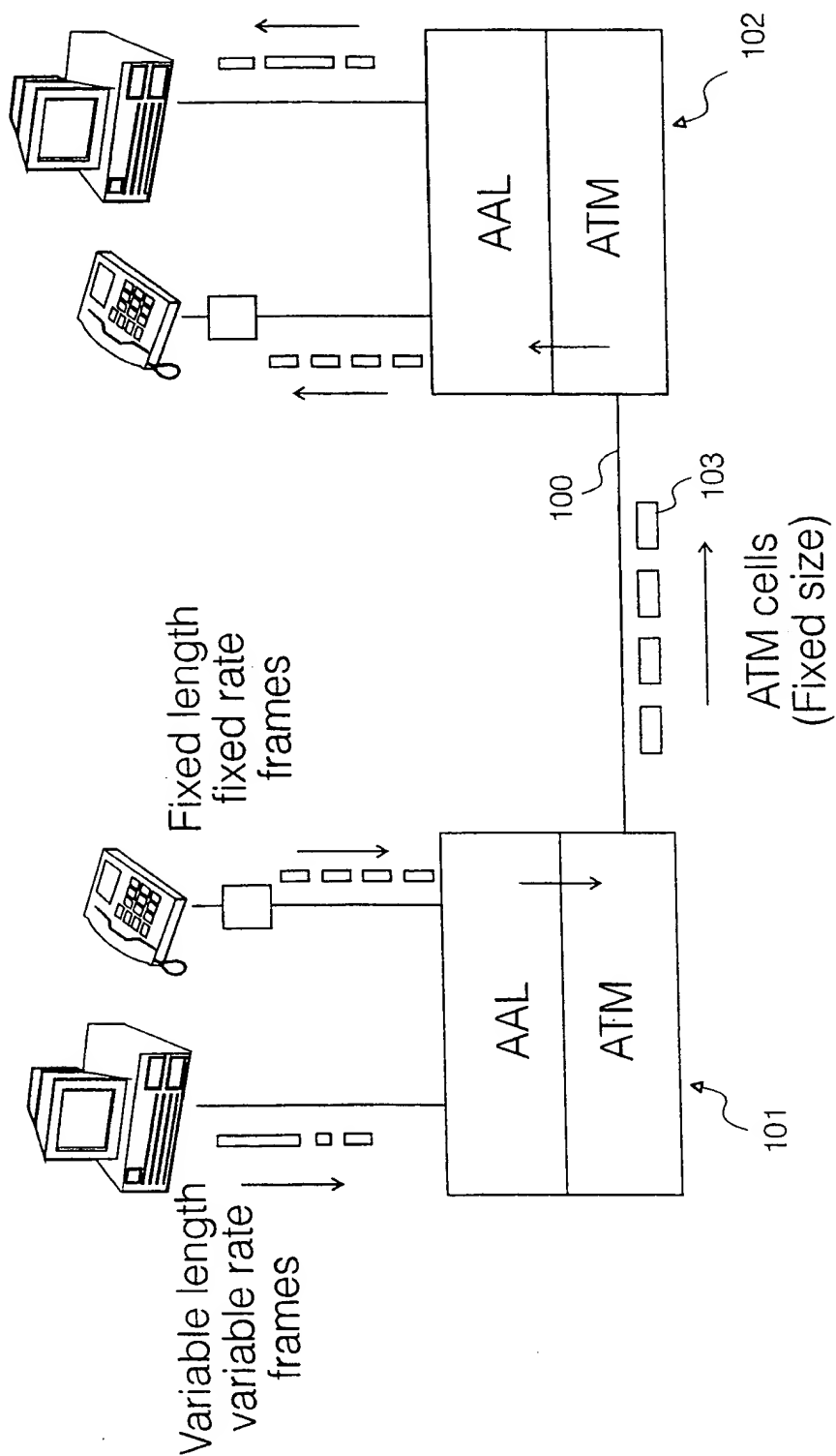
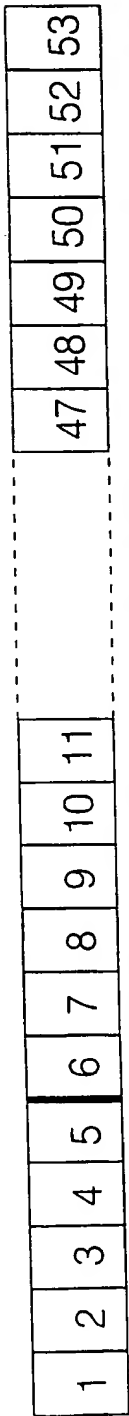


Fig. 1
(Prior Art)

ATM Cell

1 ↑ Signal
0 | level



Data Payload
(ATM-SDU)

Cell
Header

→ Octets of data bits

Fig. 2
(Prior Art)

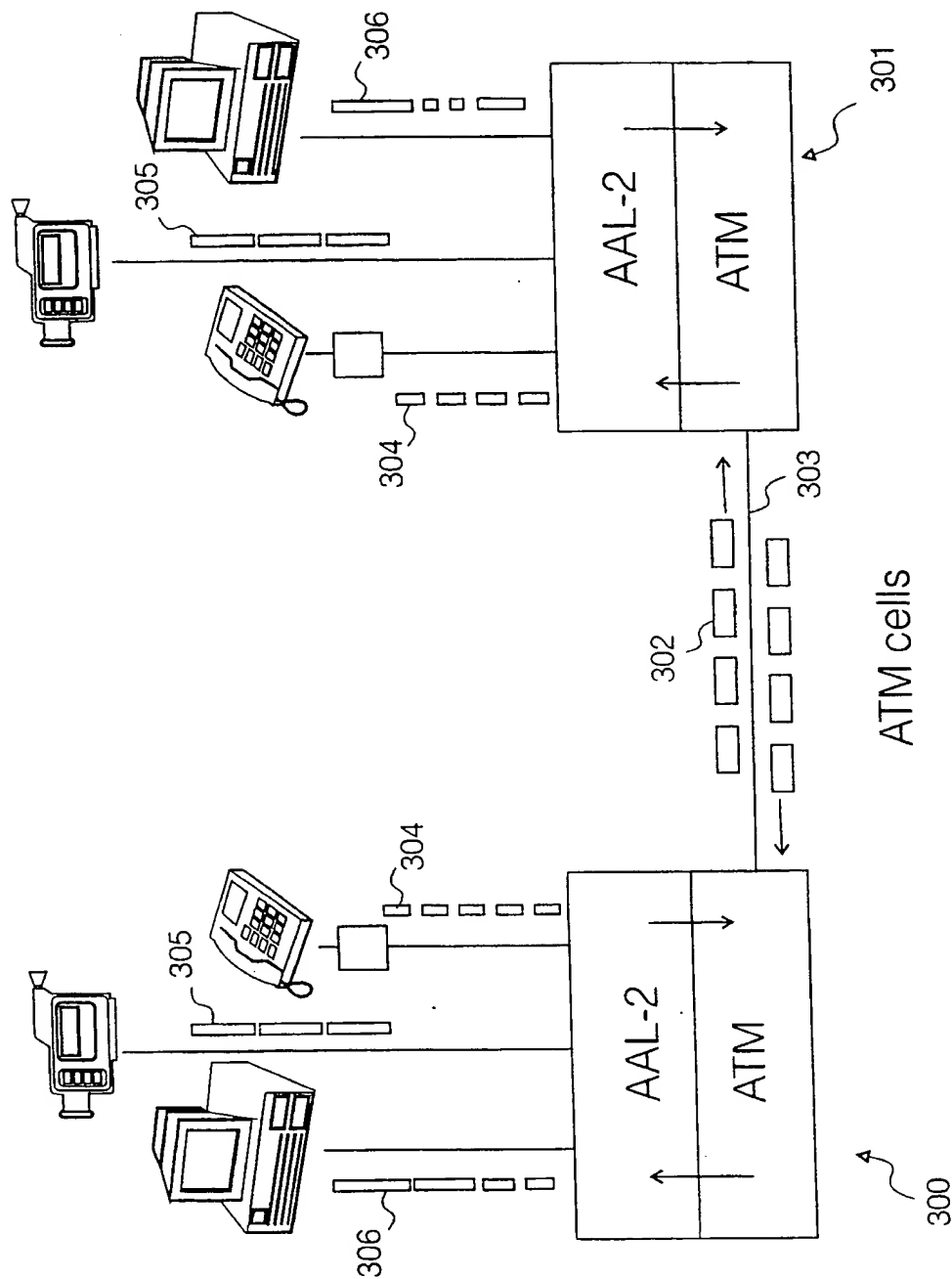


Fig. 3

4/19

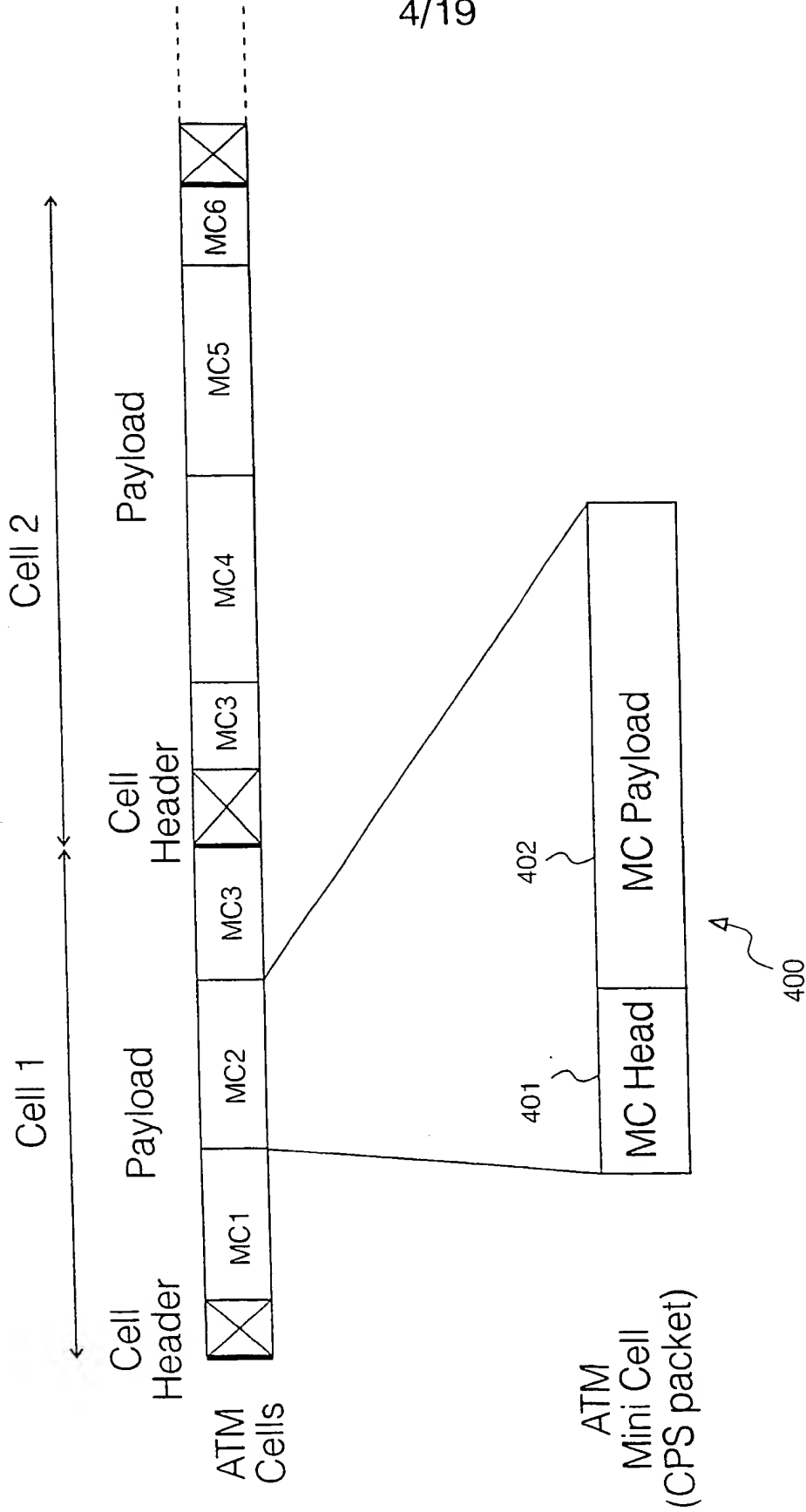


Fig. 4
(Prior Art)

5/19

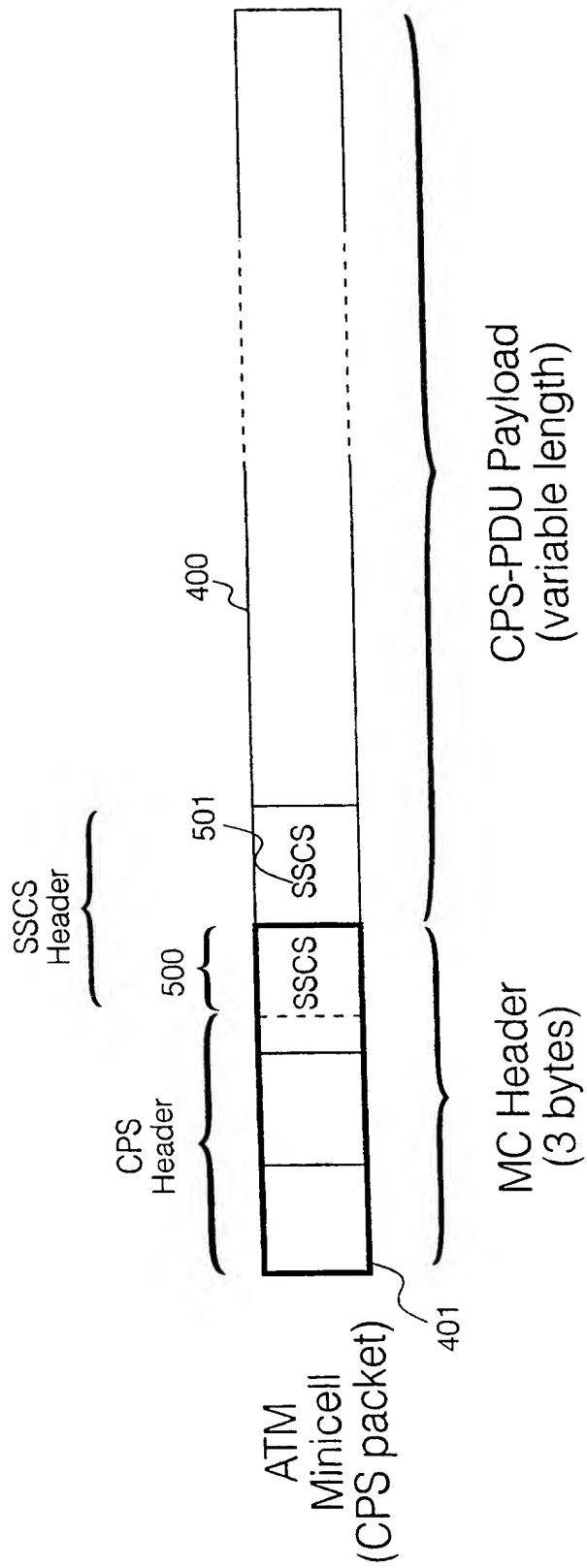
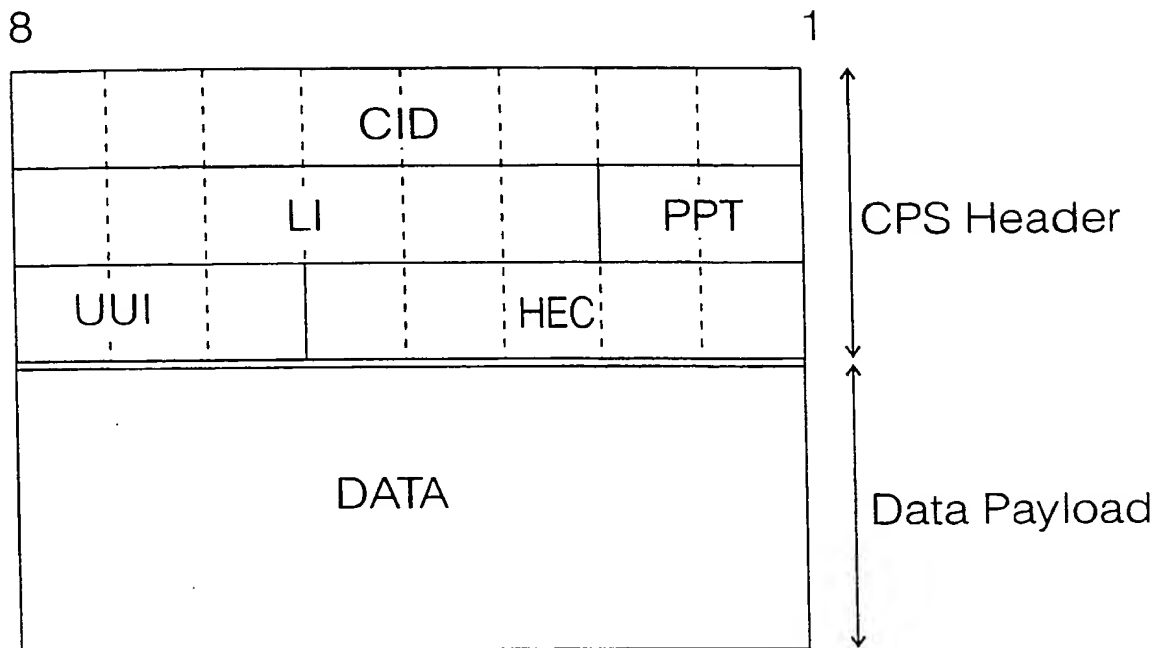


Fig. 5

6/19

ATM Minicell

CID = Channel identifier (8 bits)

LI = Length indicator (6 bits)

PPT = CPS - Packet payload type (2 bits)

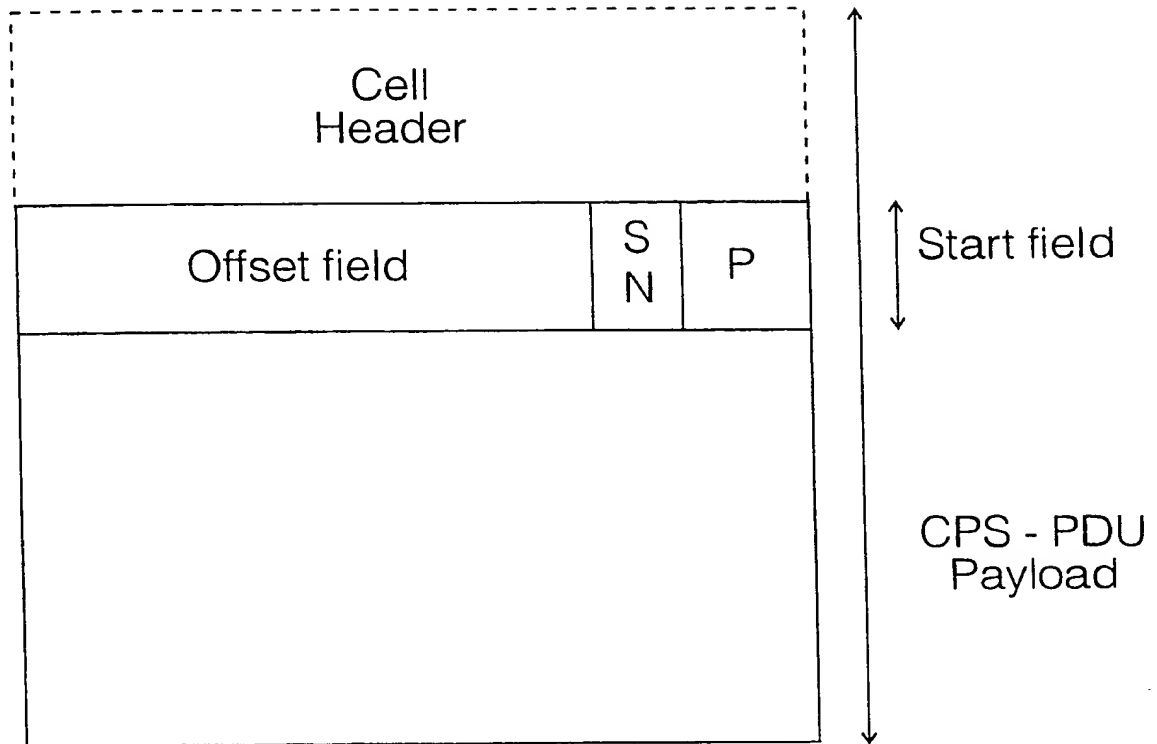
UUI = CPS - User to user indication (3 bits)

HEC = Header error control (5 bits)

DATA = User data (1 to 45 or 64 octets)

Fig. 6 (Prior Art)

7/19

ATM Minicell

OSF = Offset field (6 bits)

SN = Sequence number (1 bit)

P = Parity (1 bit)

Fig. 7 (Prior Art)

8/19

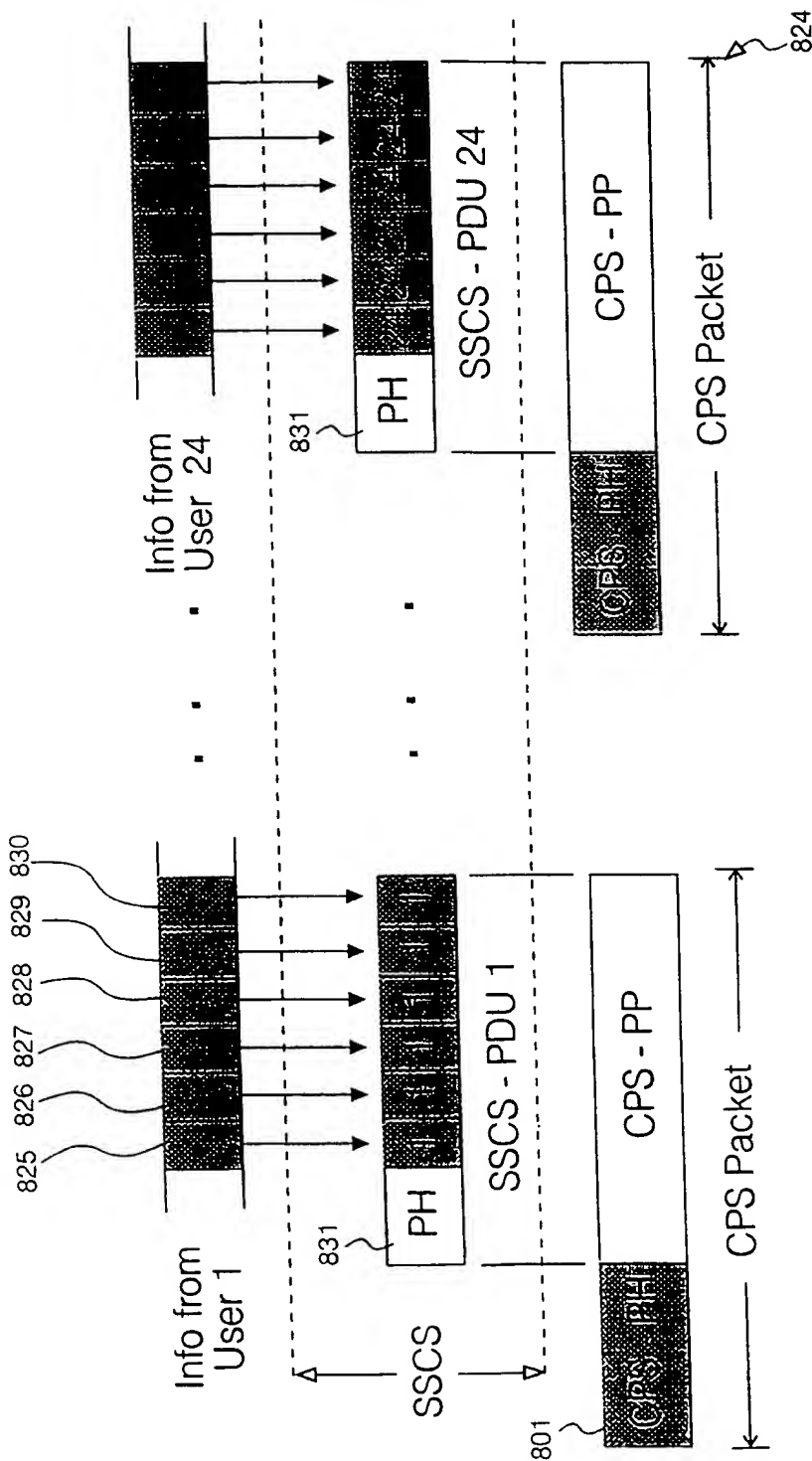


Fig. 8

9/19

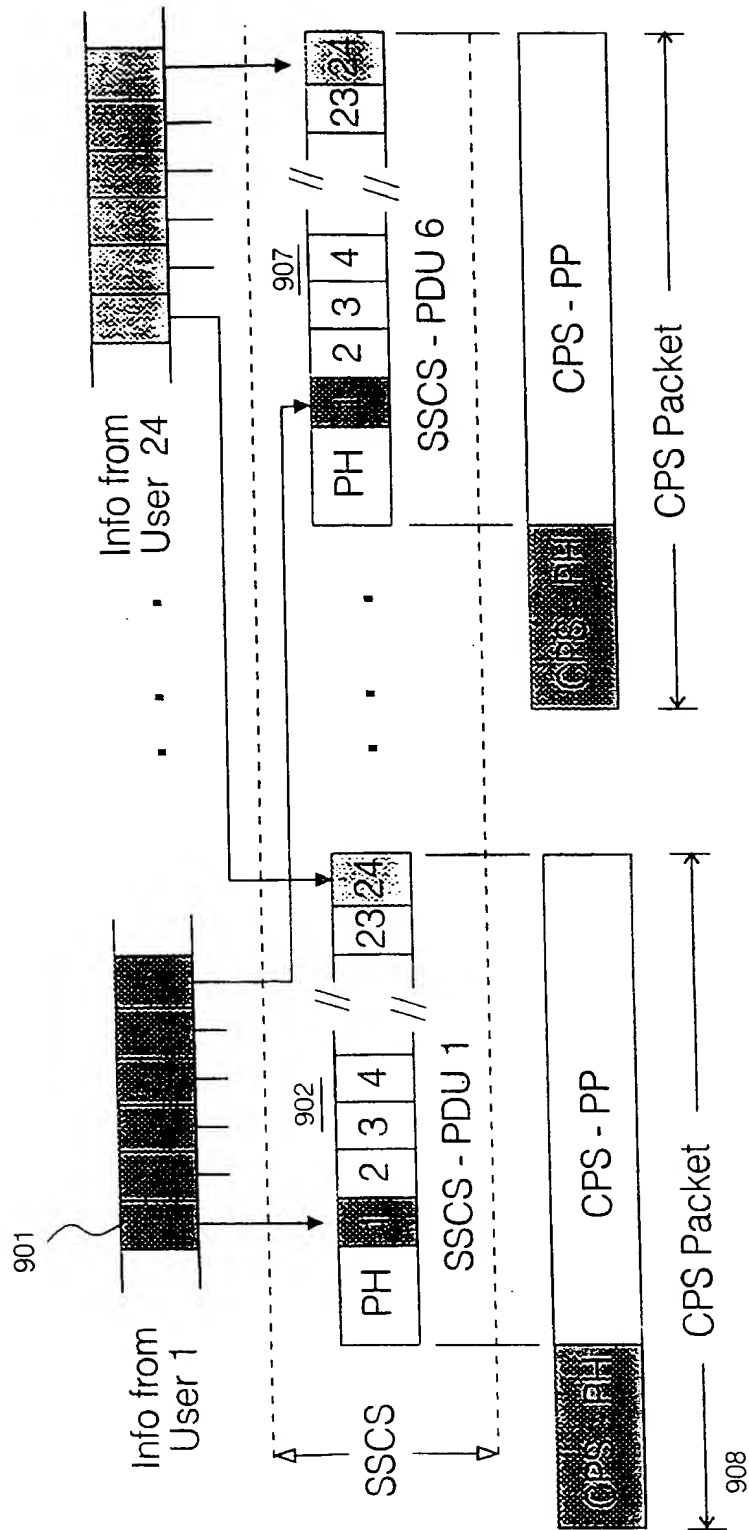


Fig. 9

10/19

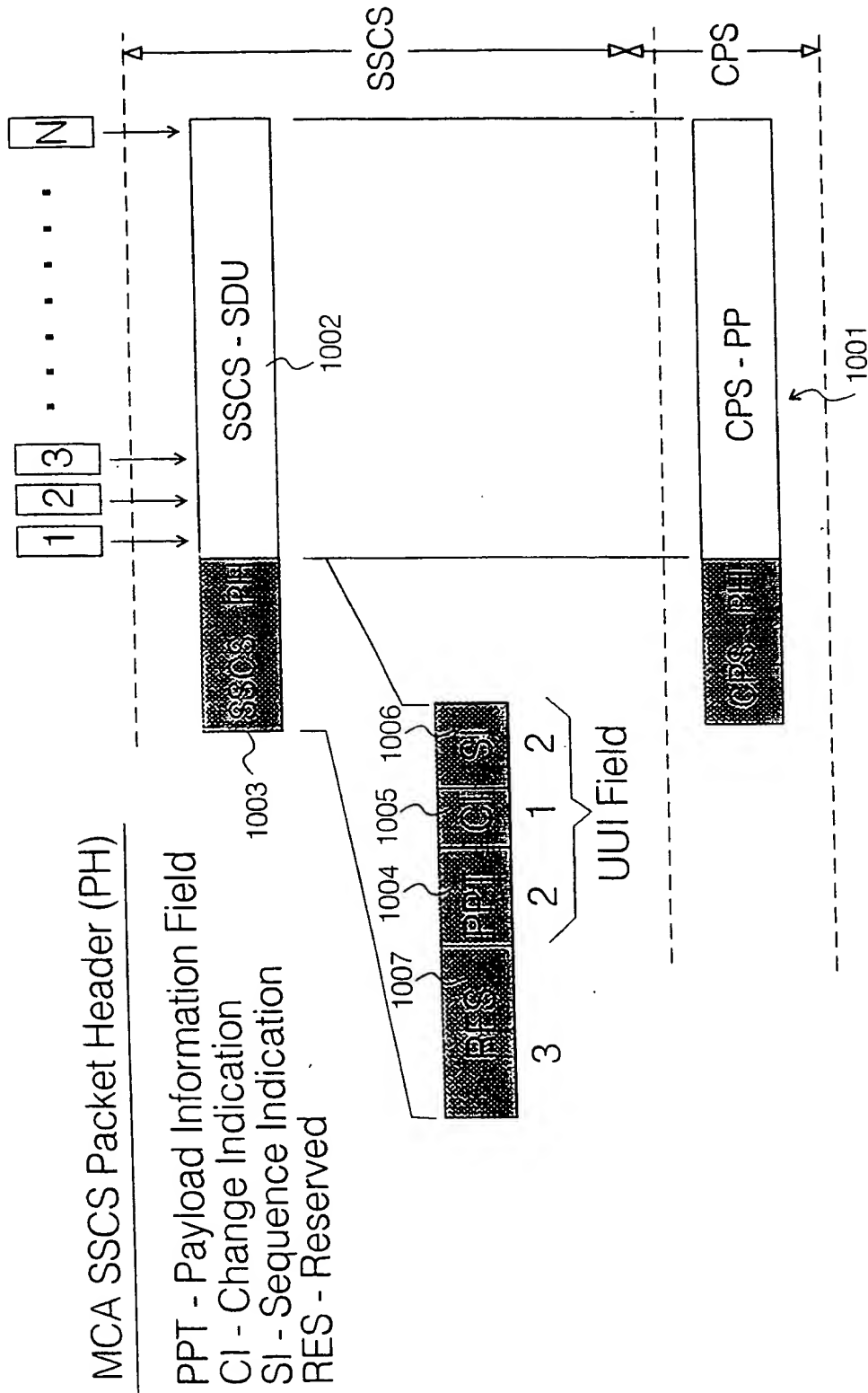


Fig. 10

11/19

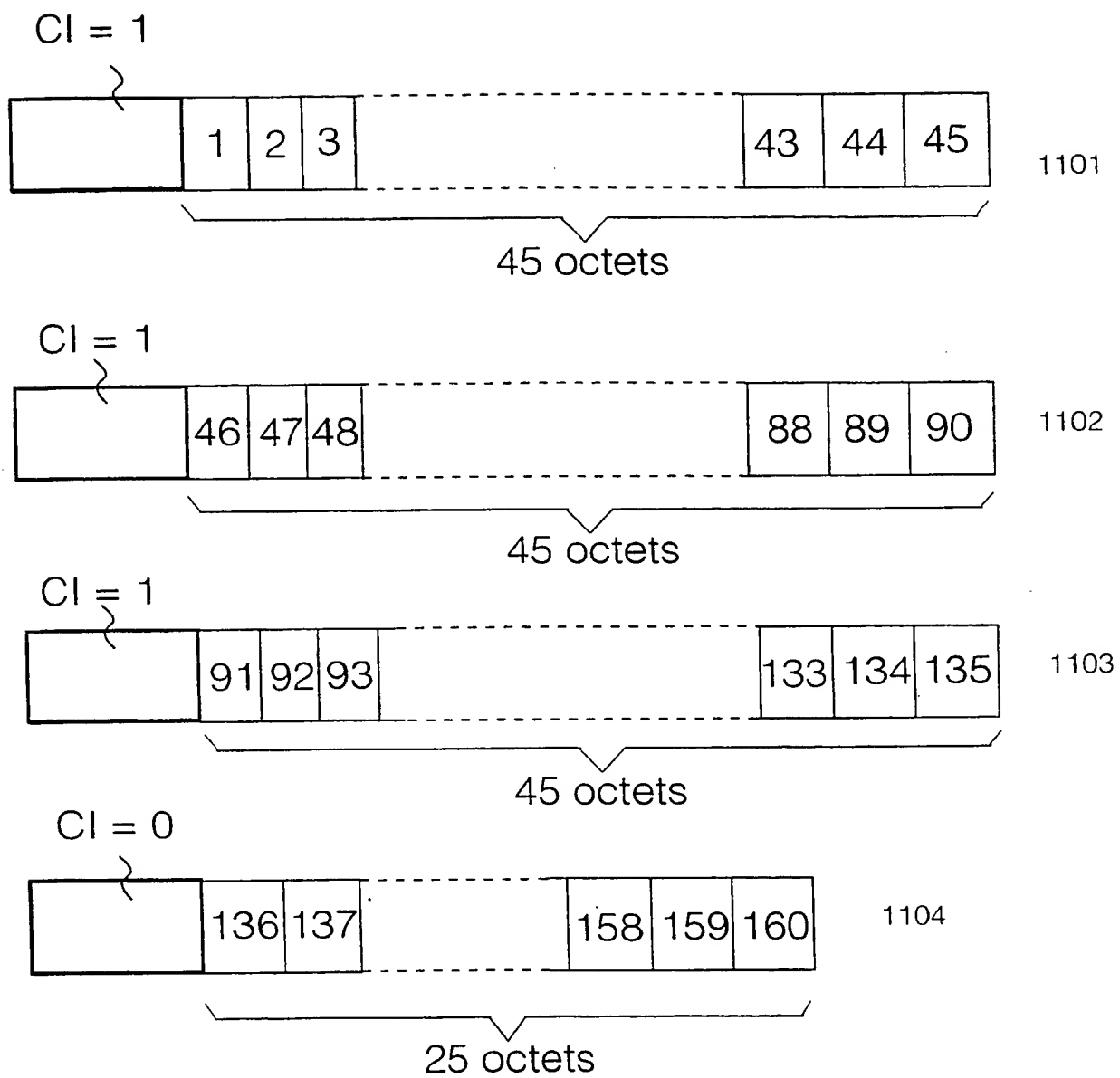


Fig. 11

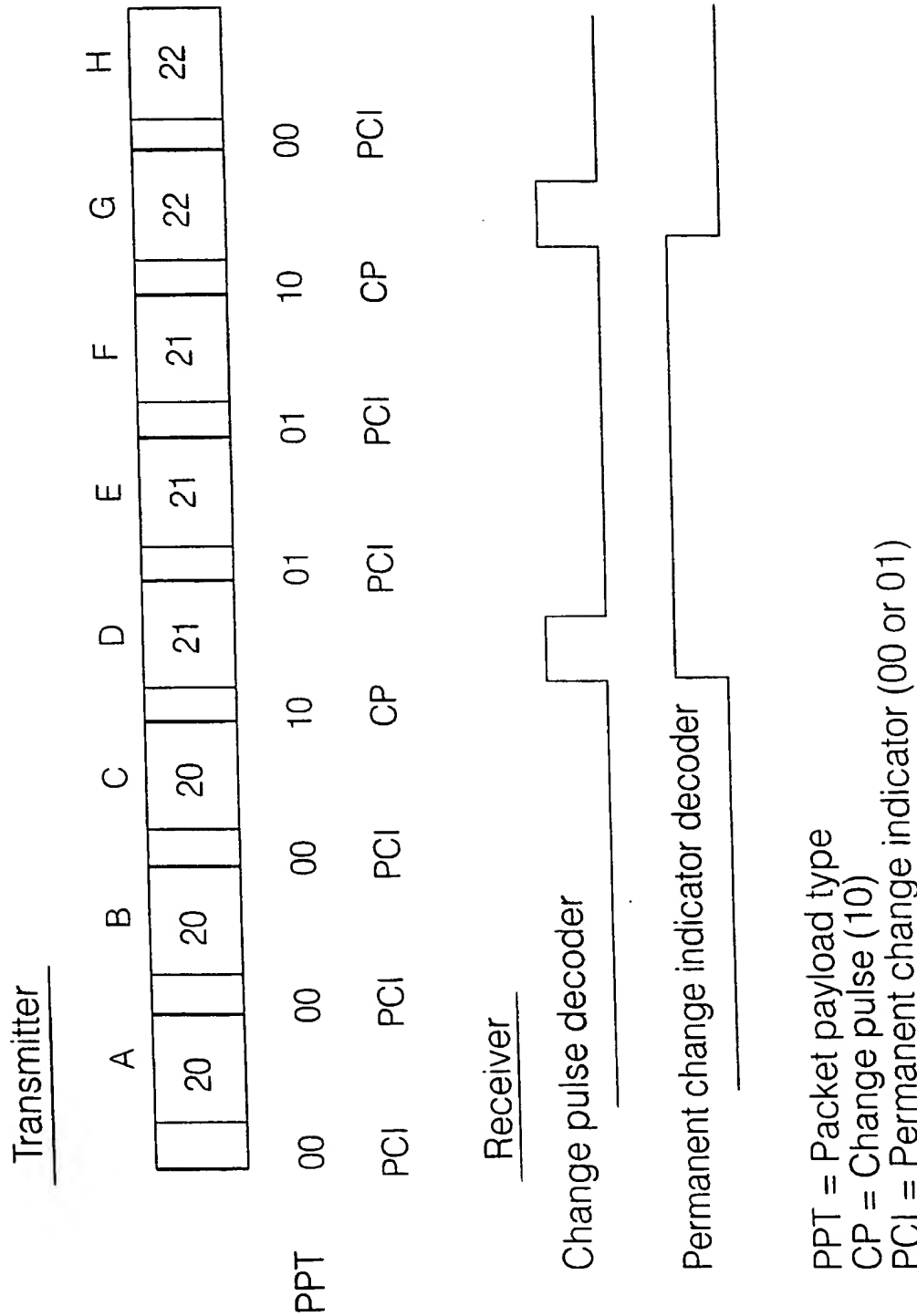


Fig. 12

13/19

TX

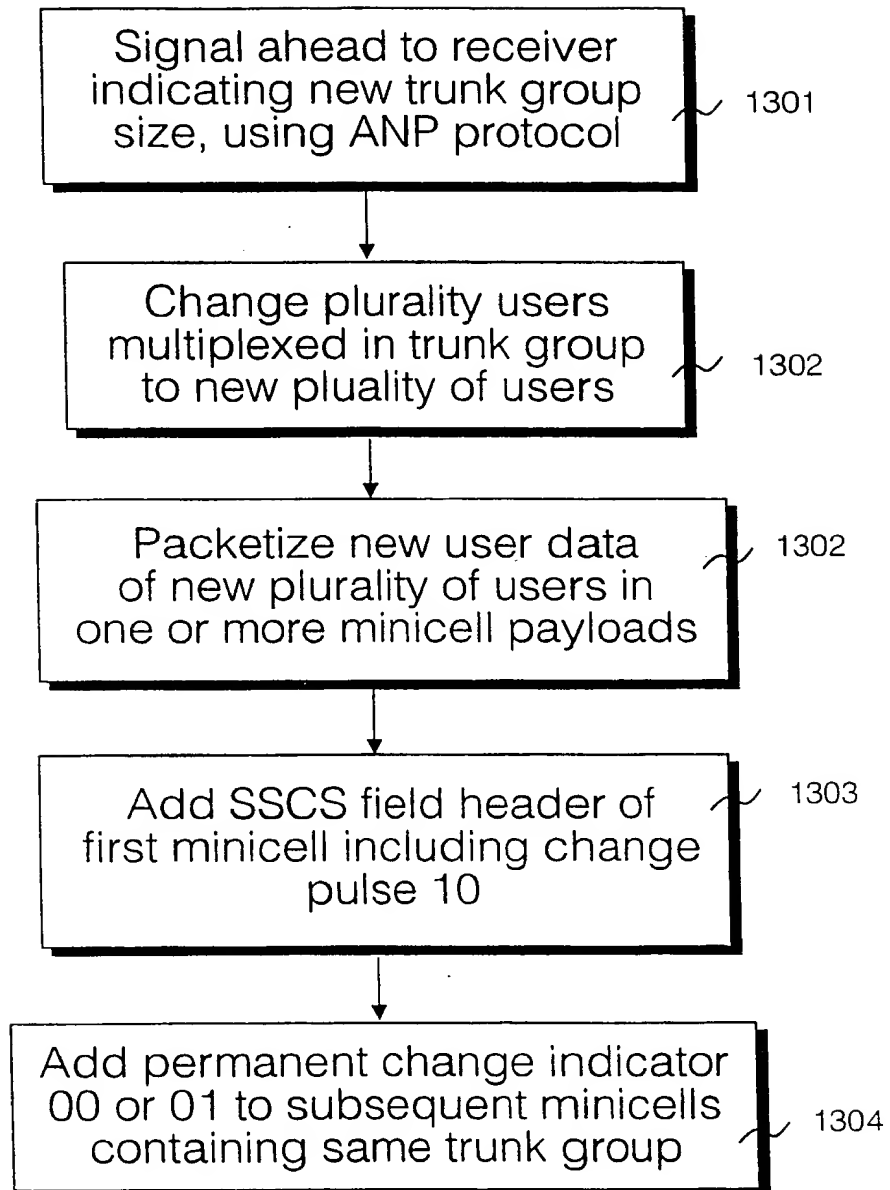


Fig. 13

14/19

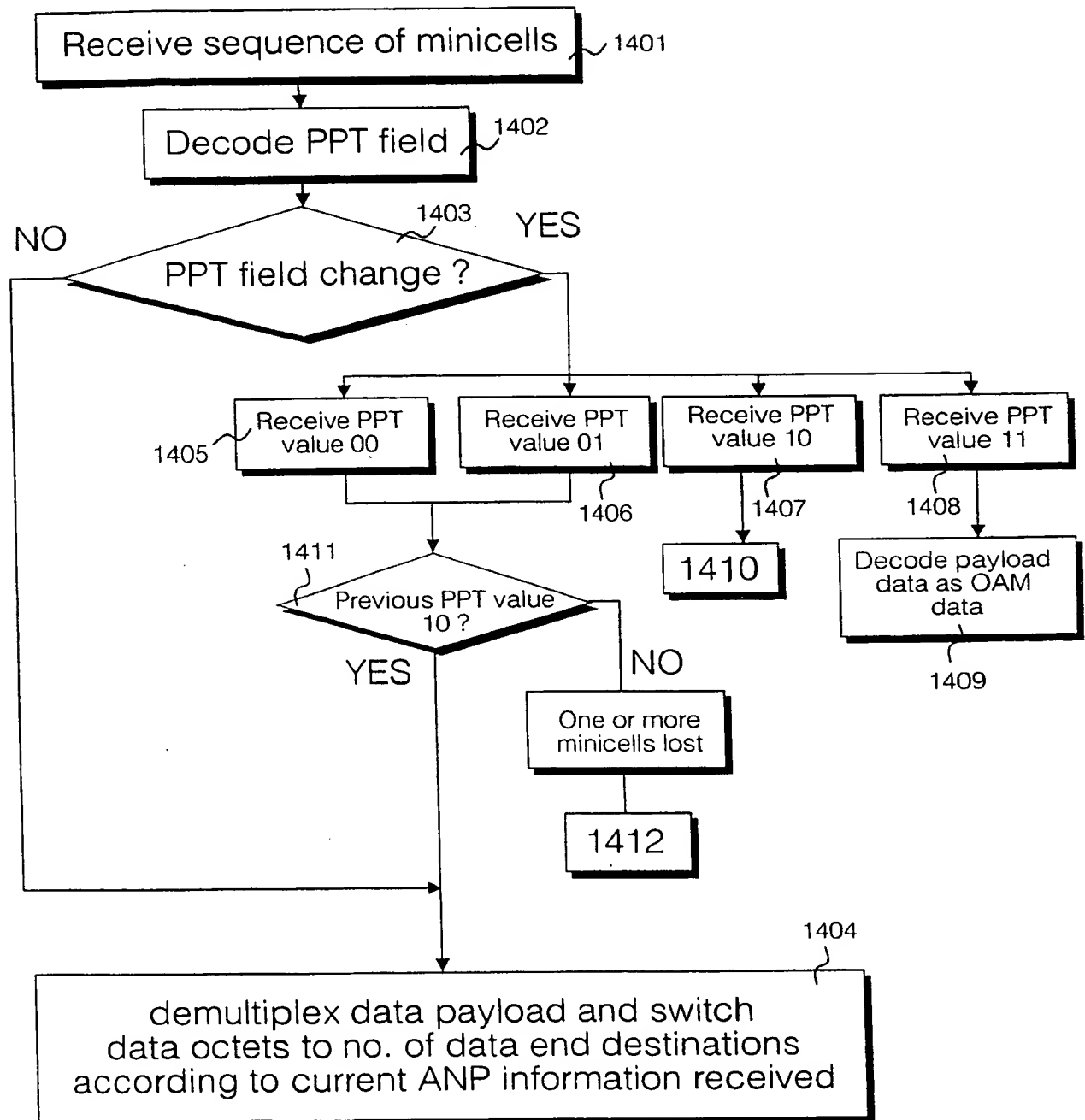
Decoder

Fig. 14

15/19

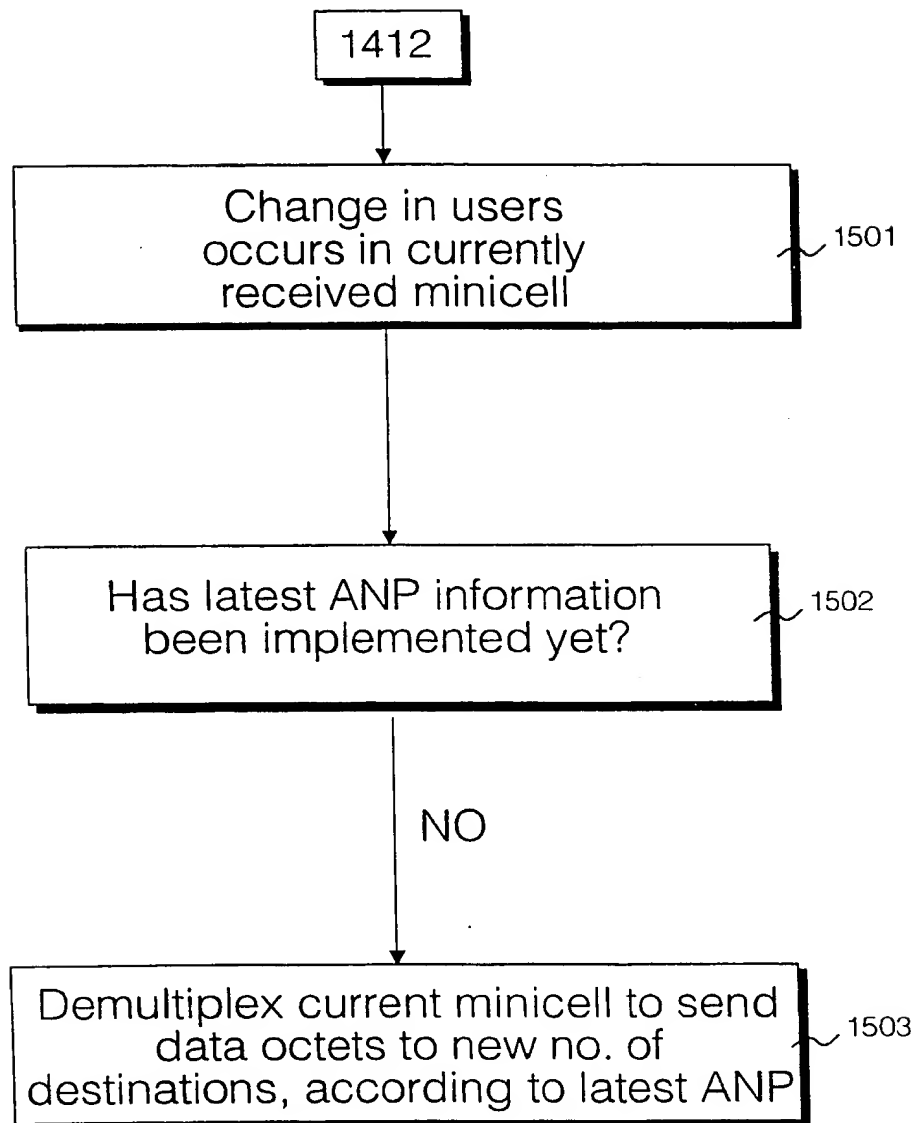
Receive PPT Value 10

Fig. 15

16/19

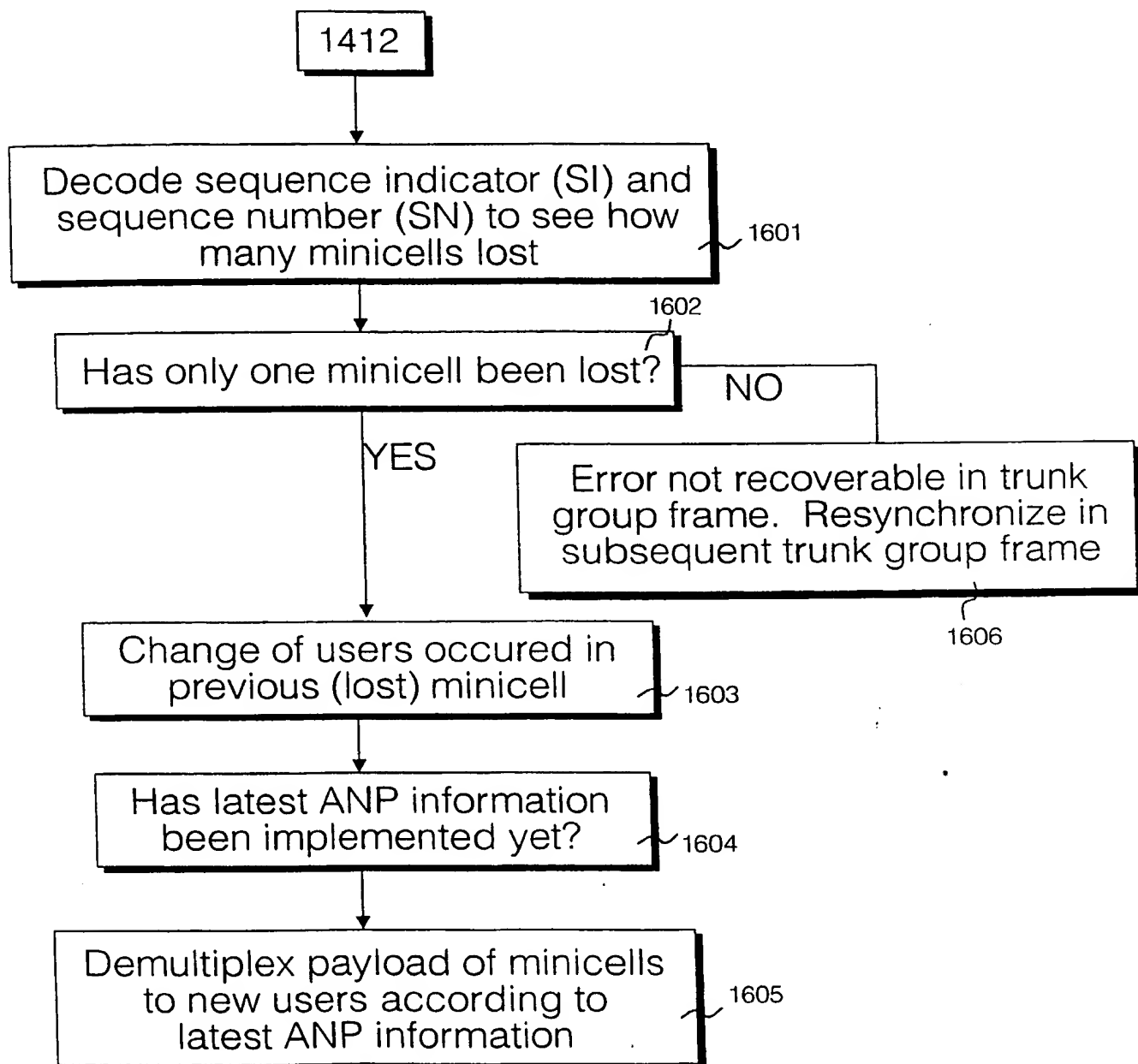
One or more minicells lost

Fig. 16

17/19

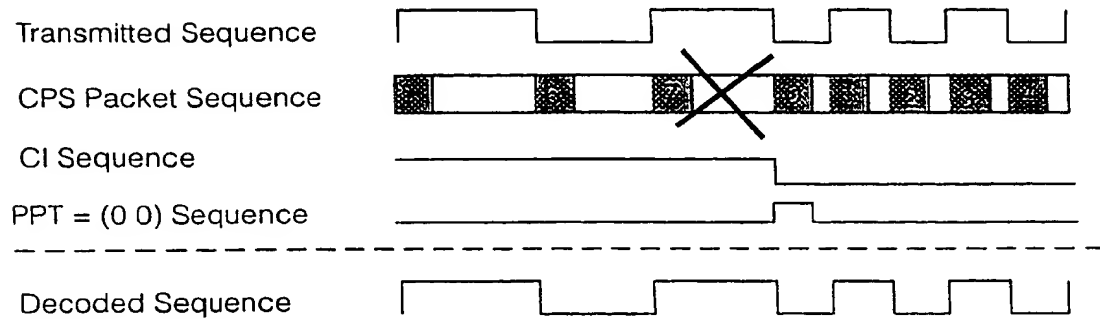
Packet lost before change implemented

Fig. 17

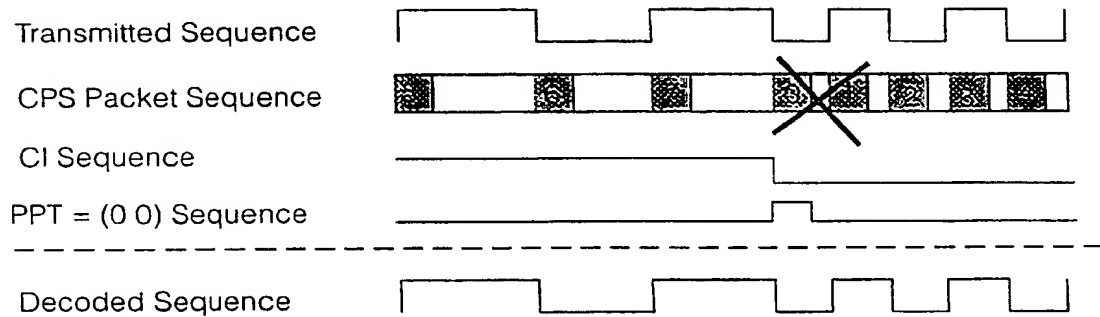
Packet lost with change indicator

Fig. 18

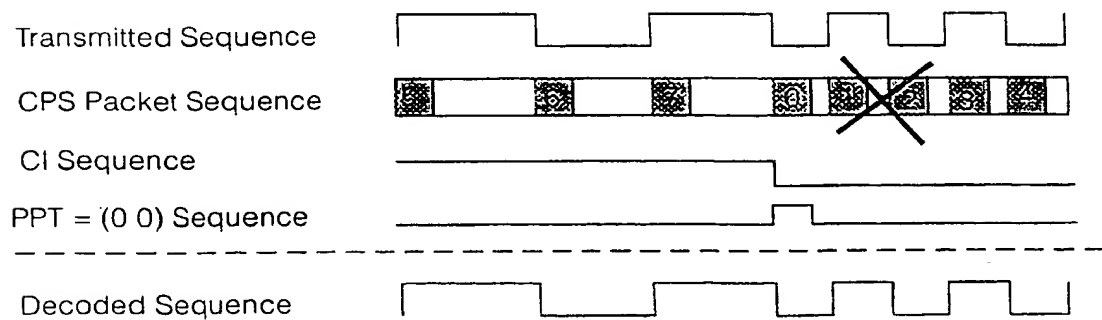
Packet lost after the change

Fig. 19

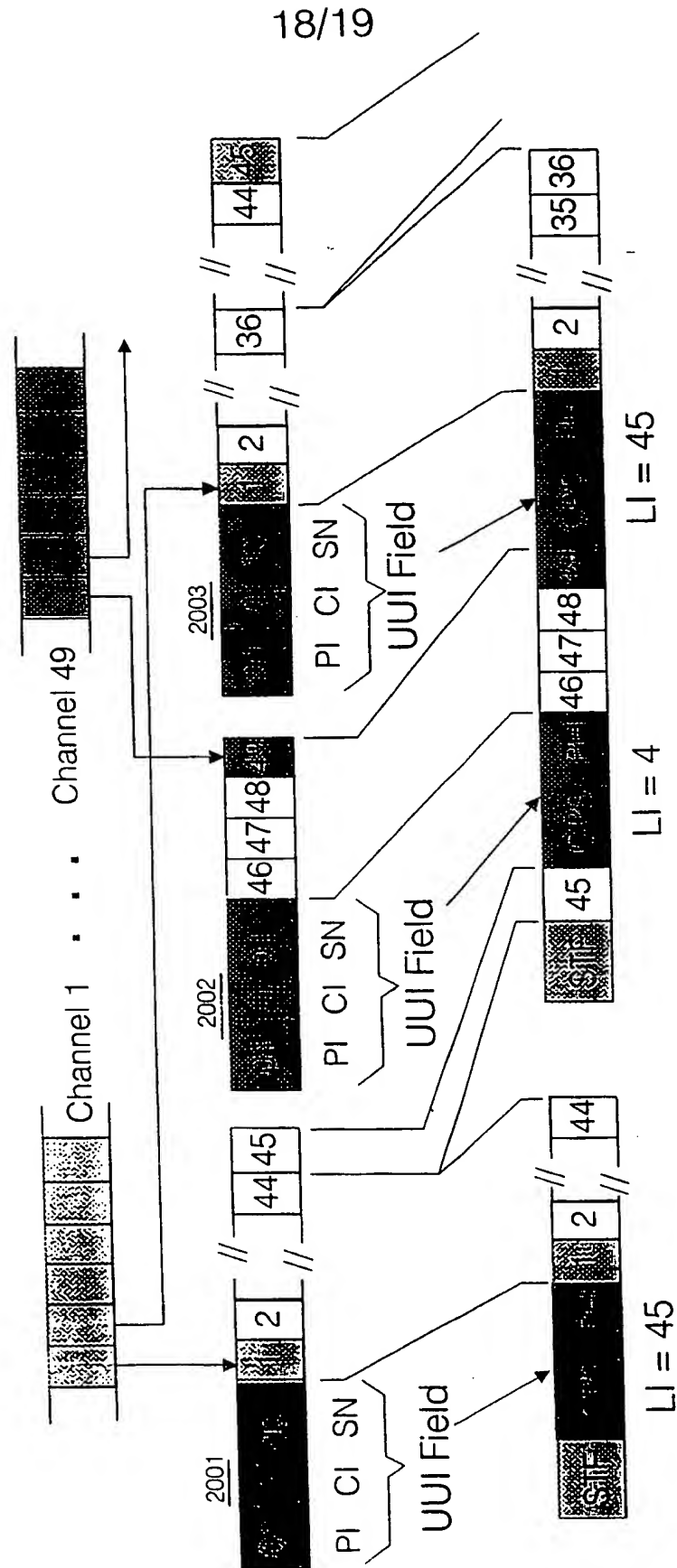


Fig. 20

19/19

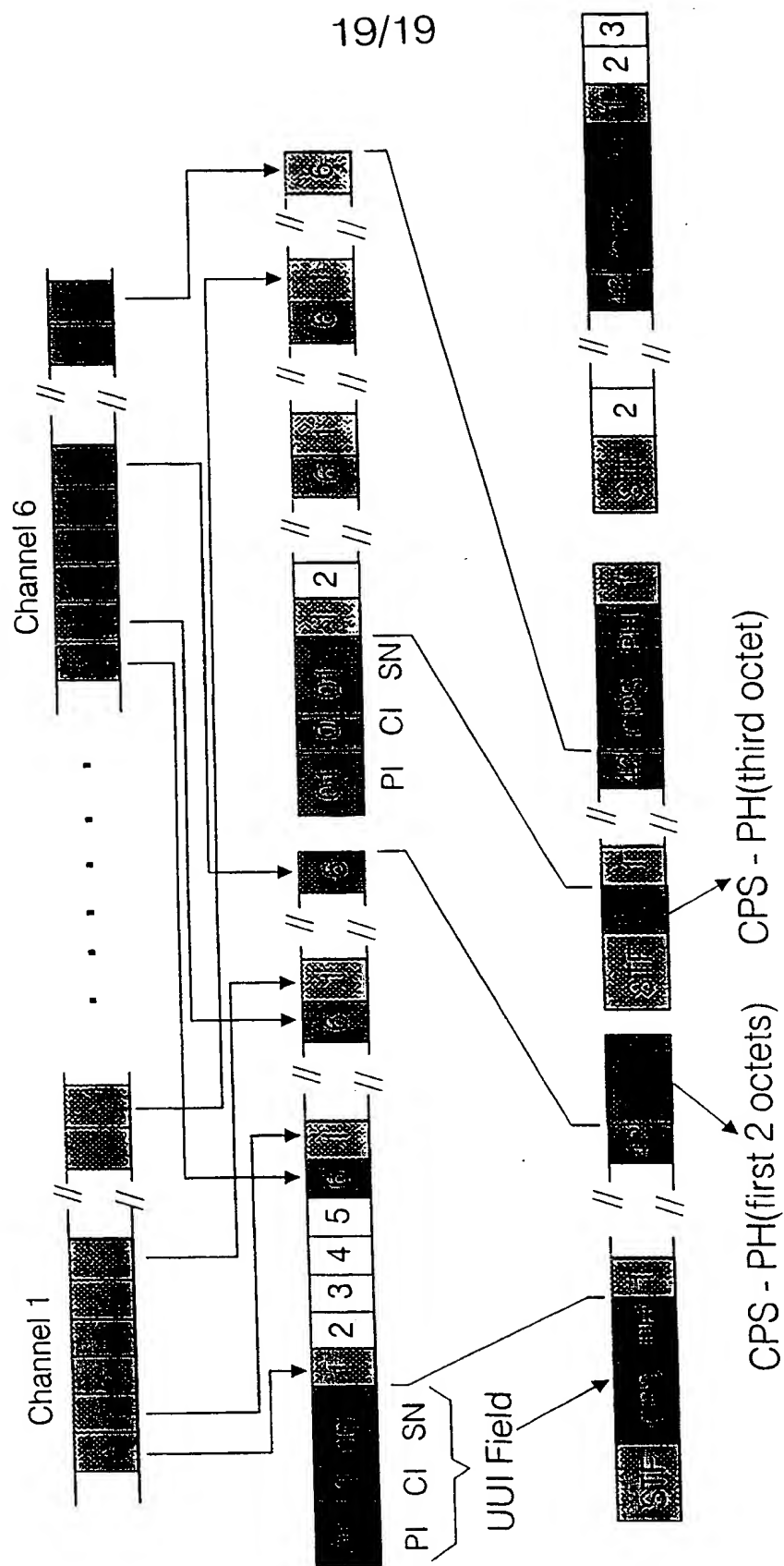


Fig. 21

INTERNATIONAL SEARCH REPORT

Intern. Application No

PCT/GB 98/00179

A. CLASSIFICATION OF SUBJECT MATTER
IPC 6 H04011/04

According to International Patent Classification (IPC) or to both national classification and IPC

B. FIELDS SEARCHED

Minimum documentation searched (classification system followed by classification symbols)

IPC 6 H040

Documentation searched other than minimum documentation to the extent that such documents are included in the fields searched

Electronic data base consulted during the international search (name of data base and, where practical, search terms used)

C. DOCUMENTS CONSIDERED TO BE RELEVANT

Category	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
X	WO 95 34977 A (ERICSSON TELEFON AB L M) 21 December 1995	1, 2, 12
Y	see page 6, line 14 - page 7, line 19 see page 11, line 31 - page 14, line 5 see page 33, line 4 - page 33, line 22; figures 24A, 24B	13, 16
X	WO 96 34478 A (ERICSSON TELEFON AB L M ; PETERSEN LARS GOERAN (SE); OLSTEDT MATS) 31 October 1996	3-7, 11, 17
Y	see page 2, line 25 - page 3, line 23 see page 7, line 24 - page 11, line 28	8-10, 13, 19, 20
	-/--	

☒ Further documents are listed in the continuation of box C

☒ Patent family members are listed in annex

Special categories of cited documents:

"A" document defining the general state of the art which is not considered to be of particular relevance

"E" earlier document but published on or after the international filing date

"L" document which may throw doubts on priority claim(s) or which is cited to establish the publication date of another citation or other special reason (as specified)

"O" document referring to an oral disclosure, use, exhibition or other means

"P" document published prior to the international filing date but later than the priority date claimed

"T" later document published after the international filing date or priority date and not in conflict with the application but cited to understand the principle or theory underlying the invention

"X" document of particular relevance: the claimed invention cannot be considered novel or cannot be considered to involve an inventive step when the document is taken alone

"Y" document of particular relevance: the claimed invention cannot be considered to involve an inventive step when the document is combined with one or more other such documents, such combination being obvious to a person skilled in the art

"Z" document member of the same patent family

Date of the actual completion of the international search

8 May 1998

Date of mailing of the international search report

18/05/1998

Name and mailing address of the ISA

European Patent Office, P.O. Box 1, Patentstrasse 1
NL-2280 HV Rijswijk
Tel: (+31-70) 340-2000; Fax: (+31-70) 340-2016

Authorized officer

Gregori, S

INTERNATIONAL SEARCH REPORT

Internat. Application No.

PCT/GB 98/00179

C.(Continuation) DOCUMENTS CONSIDERED TO BE RELEVANT

Category	Citation of document with indication where appropriate of the relevant passages	Relevant to claim No
Y	JOHNSSON M ET AL: "Support for low bitrate applications in ATM networks" IFIP WORKSHOP TC6 IFIP WORKING GROUPS 6.3 AND 6.4 FOURTH WORKSHOP ON PERFORMANCE MODELLING AND EVALUATION OF ATM NETWORKS: PARTICIPANTS PROCEEDINGS, PROCEEDINGS OF IFIP 4TH WORKSHOP ON PERFORMANCE MODELLING AND EVALUATION OF ATM NETWORKS, ILKLEY, UK., 1996, BRADFORD, UK, UNIV. BRADFORD, UK. pages 39/1-14, XP002045:06 * section 5 *	8-10, 16, 19, 20
A	-----	14, 15, 18, 21-26

INTERNATIONAL SEARCH REPORT

Information on patent family members

International Application No

PCT/GB 98/00179

Patent document cited in search report	Publication date	Patent family member(s)	Publication date
WO 9534977 A	21-12-95	SE 503317 C	13-05-96
		AU 2756295 A	05-01-96
		BR 9507985 A	12-08-97
		CA 2190459 A	21-12-95
		EP 0765557 A	02-04-97
		FI 964981 A	11-02-97
		NO 965274 A	12-02-97
		SE 9402051 A	14-12-95
WO 9634478 A	31-10-96	SE 505845 C	13-10-97
		AU 5413496 A	18-11-96
		EP 0823152 A	11-02-98
		NO 974896 A	22-12-97
		SE 9600277 A	25-10-96

THIS PAGE BLANK (USPTO)